

There are 5 problems, each with 3 parts. Each part is worth 2 points, for a total of 30 points. If any two exams are submitted with identical answers then **both** exams will receive 0 points.

**1. Induction.** For all positive integers  $n$  and  $p$  we define

$$S_p(n) := 1^p + 2^p + 3^p + \cdots + n^p.$$

In this problem you will use induction to prove that  $S_1(n) = n(n+1)/2$  for all  $n \geq 1$ .

(a) Check that the formula  $S_1(n) = n(n+1)/2$  is true for  $n = 1, 2$ , and  $3$ .

$n$	$S_1(n)$	$n(n+1)/2$	correct?
1	1	$\frac{1(1+1)}{2} = \frac{1 \cdot 2}{2} = 1$	✓
2	$1 + 2 = 3$	$\frac{2(2+1)}{2} = \frac{2 \cdot 3}{2} = 3$	✓
3	$1 + 2 + 3 = 6$	$\frac{3(3+1)}{2} = \frac{3 \cdot 4}{2} = 6$	✓

(b) Find a simple equation relating  $S_1(n)$  and  $S_1(n+1)$ .

$$S_1(n+1) = 1 + 2 + 3 + \cdots + n + (n+1)$$

$$S_1(n+1) = S_1(n) + (n+1).$$

(c) Prove that **if**  $S_1(n) = n(n+1)/2$ , **then**  $S_1(n+1) = (n+1)(n+2)/2$ . You must begin your proof with the words “Consider some  $n \geq 1$  and assume that...”.

Consider some  $n \geq 1$  and assume that  $S_1(n) = n(n+1)/2$ . In this case we have

$$\begin{aligned} S_1(n+1) &= S_1(n) + (n+1) \\ &= \frac{n(n+1)}{2} + (n+1) \\ &= (n+1) \left( \frac{n}{2} + 1 \right) \\ &= \frac{(n+1)(n+2)}{2}, \end{aligned}$$

as desired.

**2. Recurrence.** Consider the following recurrence relation:

$$\boxed{p_n = p_{n-1} + n.}$$

(a) **Assume** that  $p_0 = 3$ . In this case, make a table of  $p_n$  for  $n$  between 0 and 6.

$n$	0	1	2	3	4	5	6
$p_n$	3	4	6	9	13	18	24

(b) **Assume** that  $p_5 = 20$ . In this case, tell me the value of  $p_2$ .

We can rearrange the recurrence to get  $p_{n-1} = p_n - n$ . Then we have

$$p_4 = p_5 - 5 = 20 - 5 = 15$$

$$p_3 = p_4 - 4 = 15 - 4 = 11$$

$$p_2 = p_3 - 3 = 11 - 3 = 8.$$

(c) **Assume** that  $p_0 = 1$ . In this case, tell me a **closed formula** for  $p_n$ . [Hint: Prob 1.]

We have

$$p_0 = 1$$

$$p_1 = p_0 + 1 = 1 + 1$$

$$p_2 = p_1 + 2 = 1 + 1 + 2$$

$$p_3 = p_2 + 3 = 1 + 1 + 2 + 3$$

$$\vdots$$

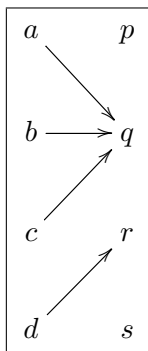
$$p_n = 1 + (1 + 2 + 3 + \cdots + n).$$

Then using the formula from Problem 1 gives  $p_n = 1 + \frac{n(n+1)}{2}$ .

**3. Functions.** Consider the sets  $X = \{a, b, c, d\}$  and  $Y = \{p, q, r, s\}$ .

(a) Draw an example of a function  $f : X \rightarrow Y$  that is **not** injective and **not** surjective.

There are 232 different correct answers. Here is one of them:

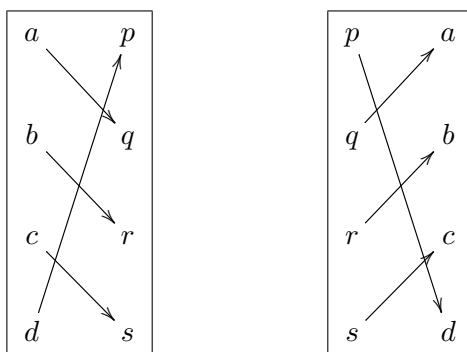


(b) Tell me the total number of different functions from  $X$  to  $Y$ .

In general the number of different functions from  $X$  to  $Y$  is  $\#Y^{\#X}$ . In this case we have  $\#X = \#Y = 4$ , hence there are  $4^4 = 256$  functions.

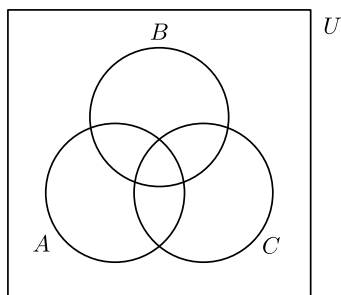
(c) Draw an example of an **invertible** function  $f : X \rightarrow Y$ . Also draw its inverse  $f^{-1}$ .

There are 24 different correct answers. Here is one of them:

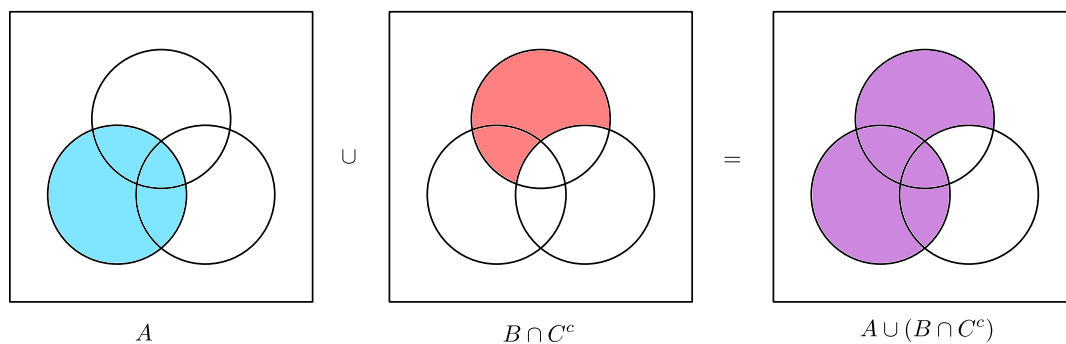


**4. Venn Diagrams.** Fix a universal set  $U$  and consider sets  $A, B, C \subseteq U$ . In this problem you will use Venn diagrams to prove that  $A \cup (B \cap C^c) = (A \cup B) \cap (A \cup C^c)$ .

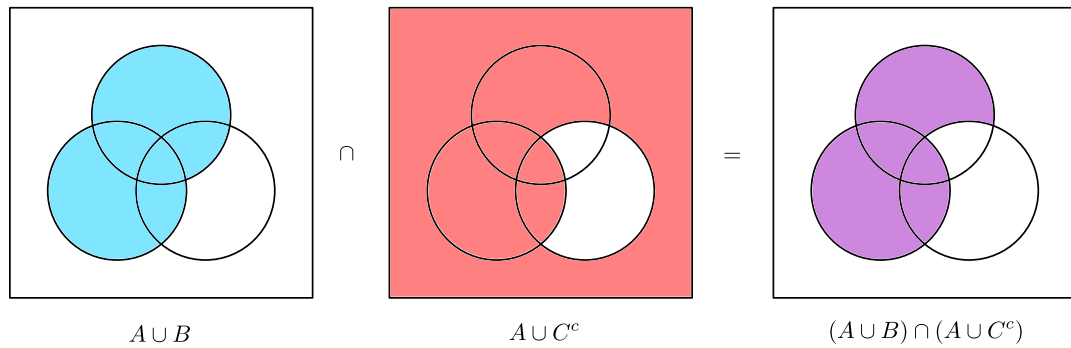
(a) Draw a general Venn diagram displaying the sets  $A$ ,  $B$ ,  $C$ , and  $U$ .



(b) Use Venn diagrams to draw the set  $A \cup (B \cap C^c)$ . Show intermediate steps.



(c) Use Venn diagrams to draw the set  $(A \cup B) \cap (A \cup C^c)$ . Show intermediate steps.



[Remark: The fact that  $A \cup (B \cap C^c) = (A \cup B) \cap (A \cup C^c)$  is an example of the distributive law.]

**5. Truth Tables.** Let  $P$  and  $Q$  be logical statements.

(a) Draw the truth table for the statement  $P \wedge Q$  (i.e.,  $P$  AND  $Q$ ).

$P$	$Q$	$P \wedge Q$
$T$	$T$	$T$
$T$	$F$	$F$
$F$	$T$	$F$
$F$	$F$	$F$

(b) Draw the truth table for the statement  $\neg(P \wedge Q)$  (i.e., NOT ( $P$  AND  $Q$ )).

$P$	$Q$	$P \wedge Q$	$\neg(P \wedge Q)$
$T$	$T$	$T$	$F$
$T$	$F$	$F$	$T$
$F$	$T$	$F$	$T$
$F$	$F$	$F$	$T$

(c) Draw the truth table for the statement  $\neg P \vee \neg Q$  (i.e., (NOT  $P$ ) OR (NOT  $Q$ )). [Hint: It may help to think about the Venn diagram of the set  $A^c \cup B^c$ .]

$P$	$Q$	$\neg P$	$\neg Q$	$\neg P \vee \neg Q$
$T$	$T$	$F$	$F$	$F$
$T$	$F$	$F$	$T$	$T$
$F$	$T$	$T$	$F$	$T$
$F$	$F$	$T$	$T$	$T$

[Remark: The fact that  $\neg(P \wedge Q) = \neg P \vee \neg Q$  is an example of de Morgan's law.]

1. Let  $n$  be a positive whole number. Find a **closed form** for the following sum:

$$1 + 2 + 3 + \cdots + n = \frac{n(n+1)}{2} = \frac{1}{2}n^2 + \frac{1}{2}n$$

2. Let  $n$  be a positive whole number. Find a **closed form** for the following sum:

$$1^2 + 2^2 + 3^2 + \cdots + n^2 = \frac{n(n+1)(2n+1)}{6} = \frac{1}{3}n^3 + \frac{1}{2}n^2 + \frac{1}{6}n$$

3. Translate the following statement into English. (Many correct answers.)

$$\forall x \in S, P(x)$$

“**Every** element  $x$  of the set  $S$  satisfies property  $P(x)$ . ”

4. Translate the following statement into English. (Many correct answers.)

$$\exists x \in S, \neg P(x)$$

“There **exists** an element  $x$  of the set  $S$  such that property  $P(x)$  does **not** hold.”

5. What is the logical relationship between the statements in Problems 3 and 4?

They are opposites. Writing symbolically, we have

$$\neg(\forall x \in S, P(x)) \equiv \exists x \in S, \neg P(x).$$

Throughout this quiz,  $P$  and  $Q$  are Boolean variables.

1. Write out the truth table for  $P \vee Q$ .

*Proof.*

$P$	$Q$	$P \vee Q$
$T$	$T$	$T$
$T$	$F$	$T$
$F$	$T$	$T$
$F$	$F$	$F$

□

2. Write out the truth table for  $P \wedge Q$ .

*Proof.*

$P$	$Q$	$P \wedge Q$
$T$	$T$	$T$
$T$	$F$	$F$
$F$	$T$	$F$
$F$	$F$	$F$

□

Now consider the Boolean function  $\varphi : \{T, F\}^2 \rightarrow \{T, F\}$  defined by the following table:

$P$	$Q$	$\varphi(P, Q)$
$T$	$T$	$T$
$T$	$F$	$F$
$F$	$T$	$T$
$F$	$F$	$T$

3. Write out the truth table for  $\neg \varphi(P, Q)$ .

*Proof.*

$P$	$Q$	$\neg \varphi(P, Q)$
$T$	$T$	$F$
$T$	$F$	$T$
$F$	$T$	$F$
$F$	$F$	$F$

□

4. Tell me a **formula** for  $\neg \varphi(P, Q)$  in terms of  $\vee, \wedge, \neg$ . [Hint: Disjunctive normal form.]

*Proof.* The only  $T$  is in the  $(T, F)$  row of the truth table. This row corresponds to the formula  $P \wedge \neg Q$  (and we can think of it as a region of the corresponding Venn diagram). Therefore the disjunctive normal form is

$$\neg \varphi(P, Q) = P \wedge \neg Q.$$

□

5. Now tell me a **formula** for  $\varphi(P, Q)$ . [Hint: Use Problem 4 and de Morgan's identity.]

*Proof.* We can read the disjunctive normal form of  $\varphi(P, Q)$  from the truth table above. There are three  $T$ 's in the rows  $(T, T)$ ,  $(F, T)$  and  $(F, F)$ . Therefore the disjunctive normal form is

$$\varphi(P, Q) = (P \wedge Q) \vee (\neg P \wedge Q) \vee (\neg P \wedge \neg Q).$$

However, this doesn't look very nice. We get a more compact formula by using Problem 4 and applying de Morgan's identity. We have

$$\neg \varphi(P, Q) = P \wedge \neg Q$$

$$\varphi(P, Q) = \neg(P \wedge \neg Q)$$

$$\varphi(P, Q) = \neg P \vee \neg \neg Q$$

$$\varphi(P, Q) = \neg P \vee Q.$$

□

[Remark: The Boolean function  $\varphi$  is very important in mathematics. In this class we will denote it by " $P \Rightarrow Q$ " :=  $\varphi(P, Q)$  and we will read the statement " $P \Rightarrow Q$ " as "if  $P$  then  $Q$ ", or " $P$  implies  $Q$ ". (WARNING: This definition might not **totally** agree with your intuition about the words "if ... then ...". For example, do you regard "if  $1 + 1 = 3$  then  $1 + 1 = 5$ " to be a true statement? I do.) Mathematical proofs and logical arguments are built from these arrows.]

1. Let  $n$  be a positive integer. Tell me a **closed formula** for the following sum:

$$\sum_{k=1}^n k = \frac{n(n+1)}{2}$$

2. Use your answer from Problem 1 to simplify the following sum:

$$\sum_{k=1}^n (2k+1) = 2 \left( \sum_{k=1}^n k \right) + \left( \sum_{k=1}^n 1 \right) = 2 \cdot \frac{n(n+1)}{2} + n = n(n+1) + n = n(n+2)$$

3. Suppose that the numbers  $p_n$  are defined by the initial condition  $p_0 = 1$  and the recurrence  $p_{n+1} = p_n + n + 2$  for all  $n \geq 0$ . Fill in the following table:

$n$	0	1	2	3	4
$p_n$	1	3	6	10	15

4. Continuing from Problem 3, use your answer from Problem 1 to find a **closed formula**:  
There are many ways to do this. Here's the least clever way:

$$\begin{aligned}
 p_0 &= 1 \\
 p_1 &= 1 + (0 + 2) \\
 p_2 &= 1 + (0 + 2) + (1 + 2) \\
 &\vdots \\
 p_n &= 1 + (0 + 2) + (1 + 2) + (3 + 2) + (4 + 2) + \cdots + ((n-1) + 2) \\
 &= 1 + \sum_{k=0}^{n-1} (k + 2) \\
 &= 1 + \left( \sum_{k=0}^{n-1} k \right) + \left( \sum_{k=0}^{n-1} 2 \right) \\
 &= 1 + \frac{(n-1)n}{2} + 2n \\
 &= \frac{2 + (n-1)n + 4n}{2} \\
 &= \frac{n^2 + 3n + 2}{2} \\
 &= \frac{(n+1)(n+2)}{2}
 \end{aligned}$$



0. (Make-up for Quiz 1) The numbers  $p_n$  are defined by the initial condition  $p_0 = 1$  and the recurrence  $p_n = p_{n-1} - n$  for all  $n \geq 1$ . Find a closed formula:

$$\begin{aligned} p_n &= 1 - 1 - 2 - 3 - \cdots - n \\ &= 1 - (1 + 2 + 3 + \cdots + n) \\ &= 1 - \frac{n(n+1)}{2} \end{aligned}$$

1. Let  $S$  be a set and for all elements  $x \in S$  let  $P(x)$  be a logical statement. Translate the following statements into English:

- “ $\forall x \in S, P(x)$ ”

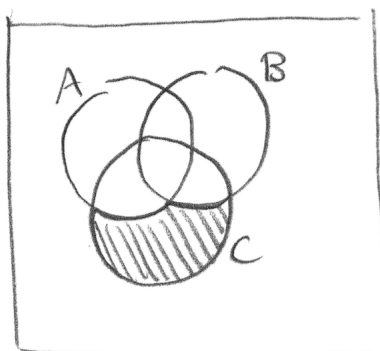
“For all elements  $x$  in  $S$ , the statement  $P(x)$  is true.”  
or “The statement  $P(x)$  holds for every element  $x$  in  $S$ .”  
or “Every element  $x$  of  $S$  satisfies  $P(x)$ .”

- “ $\exists x \in S, \neg P(x)$ ”

“There exists an element  $x$  in  $S$  such that  $P(x)$  is false.”

2. Let  $A, B, C$  be subsets of the universal set  $U$ . Use a Venn diagram to illustrate the set

$$(A \cup B)^c \cap C.$$



3. Complete the following truth table:

$P$	$Q$	$P \wedge Q$	$P \vee Q$	$P \Rightarrow Q$
$T$	$T$	$T$	$T$	$T$
$T$	$F$	$F$	$T$	$F$
$F$	$T$	$F$	$T$	$T$
$F$	$F$	$F$	$F$	$T$

4. Consider the Boolean function  $\varphi(P, Q)$  defined as follows:

$P$	$Q$	$\varphi(P, Q)$
$T$	$T$	$F$
$T$	$F$	$T$
$F$	$T$	$F$
$F$	$F$	$T$

Write down two different algebraic formulas for this function:

(There are infinitely many correct answers.)

$$\varphi(P, Q) = (P \wedge \neg Q) \vee (\neg P \wedge \neg Q)$$

$$\varphi(P, Q) = \neg Q$$

**Problem 1.**

- (a) Draw Pascal's Triangle down to the 7th row.

$$\begin{array}{ccccccccccc}
 & & & & & & 1 & & & & & \\
 & & & & & 1 & & 1 & & & & \\
 & & & 1 & & 2 & & 1 & & & & \\
 & & 1 & & 3 & & 3 & & 1 & & & \\
 & 1 & & 4 & & 6 & & 4 & & 1 & & \\
 1 & & 1 & & 5 & & 10 & & 10 & & 5 & & 1 \\
 & 1 & & 6 & & 15 & & 20 & & 15 & & 6 & & 1 \\
 & & 1 & & 7 & & 21 & & 35 & & 35 & & 21 & & 7 & & 1
 \end{array}$$

- (b) Use the triangle to expand  $(1 + x)^7$ .

$$(1 + x)^7 = 1 + 7x + 21x^2 + 35x^3 + 35x^4 + 21x^5 + 7x^6 + x^7$$

- (c) Use the triangle to evaluate the following sum:

$$\begin{aligned}
 \sum_{k=0}^4 (-1)^k \binom{7}{k} &= \binom{7}{0} - \binom{7}{1} + \binom{7}{2} - \binom{7}{3} + \binom{7}{4} \\
 &= 1 - 7 + 21 - 35 + 35 \\
 &= 15
 \end{aligned}$$

**Problem 2.** Let the sequence  $S_0, S_1, S_2, \dots$  be defined by the following initial condition and recurrence relation:

$$S_n := \begin{cases} 1 & \text{if } n = 0, \\ S_{n-1} + 2^{n-1} & \text{if } n \geq 1. \end{cases}$$

(a) Fill in the following table:

$n$	0	1	2	3	4	5
$S_n$	1	2	4	8	16	32

(b) Try to guess a simple formula for  $S_n$ .

I guess that  $S_n = 2^n$  for all  $n \geq 0$ .

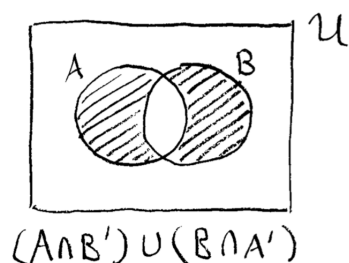
(c) Use induction to prove that your formula is correct.

- *Base Case.* If  $n = 0$  then we have  $S_0 = 1 = 2^0$ . ✓
- *Induction Step.* Now fix some  $n \geq 0$  and assume for induction that  $S_n = 2^n$ . In this case we want to prove that  $S_{n+1} = 2^{n+1}$ . Indeed, we observe that

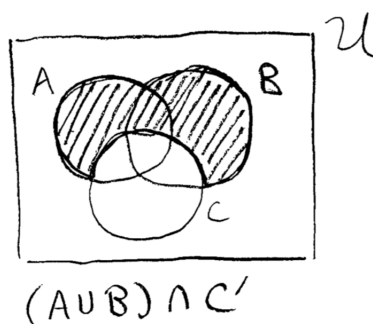
$$\begin{aligned} S_{n+1} &= S_n + 2^n && \text{by definition} \\ &= 2^n + 2^n && \text{by assumption} \\ &= 2 \cdot 2^n \\ &= 2^{n+1}. \end{aligned}$$

1. **Venn Diagrams.** Let  $A, B, C \subseteq U$  be subsets of the universal set.

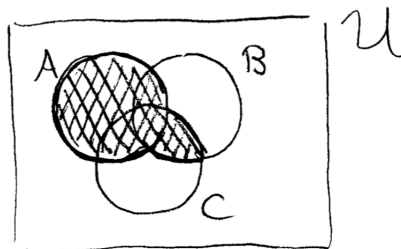
(a) Draw a Venn diagram to show the set  $(A \cap B') \cup (B \cap A')$ .



(b) Draw a Venn diagram to show the set  $(A \cup B) \cap C'$ .



(c) Tell the name of the following set (many correct answers):



The shortest name of the set is  $A \cup (B \cap C)$ . The disjunctive normal form is  $(A \cap B \cap C) \cup (A \cap B' \cap C) \cup (A \cap B \cap C') \cup (A \cap B' \cap C') \cup (A' \cap B \cap C)$ .

Or we can compute the disjunctive normal form of the unshaded region. Then we take the complement and apply de Morgan's law:

$$\begin{aligned} & [(A' \cap B' \cap C) \cup (A' \cap B \cap C') \cup (A' \cap B' \cap C')] ' \\ &= (A' \cap B' \cap C)' \cap (A' \cap B \cap C')' \cap (A' \cap B' \cap C')' \\ &= (A \cup B \cup C') \cap (A \cup B' \cup C) \cap (A \cup B \cup C). \end{aligned}$$

The resulting expression is called the *conjunctive normal form*.

**2. Boolean Functions.** Consider the Cartesian product set:

$$\{T, F\}^n := \underbrace{\{T, F\} \times \{T, F\} \times \cdots \times \{T, F\}}_{n \text{ times}}.$$

(a) Tell me the number of elements of the set  $\{T, F\}^n$ .

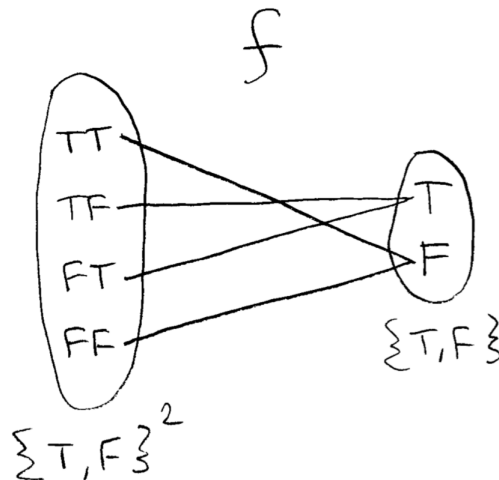
$$\#\{T, F\}^n = \underbrace{\#\{T, F\} \times \#\{T, F\} \times \cdots \times \#\{T, F\}}_{n \text{ times}} = \underbrace{2 \times 2 \times \cdots \times 2}_{n \text{ times}} = 2^n$$

(b) Count the functions  $f : \{T, F\}^2 \rightarrow \{T, F\}$  with two inputs and one output.

The number of functions from  $A$  and  $B$  is  $(\#B)^{\#A}$ . Hence the number of functions from  $\{T, F\}^2$  to  $\{T, F\}$  is

$$(\#\{T, F\})^{(\#\{T, F\}^2)} = 2^{(2^2)} = 2^4 = 16.$$

(c) Fill in the arrows for the function  $f(P, Q) = (P \wedge \neg Q) \vee (Q \wedge \neg P)$ :



Remark: This is just one of the 16 Boolean functions with two inputs and one output. It is commonly called XOR. This function also appeared in Problem 1(a).