**Problem 1. Generalized Binomial Coefficients.** Let C(n,k) denote the coefficient of  $a^k b^{n-k}$  in the expansion of  $(a+b)^n$ . We proved in class that

$$C(n,k) = \frac{n!}{k!(n-k)!}$$

for all relevant values of n and k.

(a) It seems that the above formula only makes sense when  $n, k \in \mathbb{Z}$  with  $0 \le k \le n$ . However, if we rewrite the formula as

$$C(n,k) = \frac{n!}{k!(n-k)!} = \frac{(n)_k}{k!}$$

where  $(n)_k := n(n-1)(n-2)\cdots(n-(k-1))$ , then we can define C(z,k) for any integer  $k \ge 0$  and for **any number** z whatsoever. Use this definition to prove that

$$C(-n,k) = (-1)^k \cdot C(n+k-1,k)$$

for all integers  $n, k \in \mathbb{Z}$  with  $k \ge 0$ .

(b) If z is any number and k is a negative integer then we will define C(z,k) = 0. Use induction and the result from part (a) to prove that for all integers  $n, k \in \mathbb{Z}$  (even for negative integers) we have

$$C(n,k) = C(n-1,k-1) + C(n-1,k).$$

[Hint: You already know this is true for  $n \ge 0$ . Use induction to prove it for n < 0.]

## Problem 2. Generalization of Fermat's little Theorem.

- (a) Let  $a, b, c \in \mathbb{Z}$  with gcd(a, b) = 1. If a|c and b|c, prove that ab|c. [Hint: Use Bézout to write ax + by = 1 for some  $x, y \in \mathbb{Z}$  and multiply both sides by c.]
- (b) The RSA cryptosystem is based on the following generalization of Fermat's little Theorem: For all integers  $a, p, q \in \mathbb{Z}$  with  $p \neq q$  prime and gcd(a, pq) = 1 we have

$$[a^{(p-1)(q-1)}]_{pq} = [1]_{pq}.$$

Prove this. [Hint: The condition gcd(a, pq) = 1 implies that  $p \nmid a$  and  $q \nmid a$ . We want to show that pq divides  $a^{(p-1)(q-1)} - 1$ . First observe that q does not divide  $a^{p-1}$  otherwise Euclid's Lemma implies that q|a. Then Fermat's little Theorem implies that q divides  $(a^{p-1})^{q-1} - 1 = a^{(p-1)(q-1)} - 1$ . A parallel argument shows that p divides  $a^{(p-1)(q-1)} - 1$ . Now use part (a).]

**Problem 3. RSA Cryptosystem.** You set up an RSA cipher with public key (23, 55) and private key (7, 55). I sent you the following message using the numbers 1-26 for letters of the alphabet, 27 for period, 28 for space, and 29 for exclamation point:

 $\begin{matrix} [25, 17, 1, 49, 11, 39, 7, 51, 20, 2, 7, 23, 15, 1, 2, 49, 14, 49, 13, 7, 1, 8, 20, 21, \\ 25, 7, 1, 8, 39, 25, 2, 1, 27, 25, 7, 52, 1, 25, 17, 15, 52, 1, 25, 14, 27, 39, 48, 17, \\ 1, 33, 15, 7, 1, 7, 13, 2, 15, 1, 25, 7, 12, 14, 49, 25, 15, 2, 7, 8, 2, 15, 1, 11, 24 \end{matrix}$ 

Decrypt the message.