

**Problem 1.** Let  $X$  and  $Y$  be **finite** sets.

- (a) If there exists a **surjective** function  $f : X \rightarrow Y$ , prove that  $|X| \geq |Y|$ .
- (b) If there exists an **injective** function  $g : X \rightarrow Y$ , prove that  $|X| \leq |Y|$ .
- (c) If there exists a **bijective** function  $h : X \rightarrow Y$ , prove that  $|X| = |Y|$ .

[Hint: For parts (a) and (b), for each  $y \in Y$  let  $d(y)$  be the number of arrows pointing to  $y \in Y$ . What happens if you sum the numbers  $d(y)$  for all  $y \in Y$ ? Recall the **definitions** from the course notes.]

*Proof.* Let  $X, Y$  be sets and consider a function  $f : X \rightarrow Y$ . Let  $d(y)$  denote the number of arrows of  $f$  pointing to  $y \in Y$  (this is the same as the number of  $x \in X$  such that  $f(x) = y$ ). If we sum the numbers  $d(y)$  over  $y \in Y$  we get the total number of arrows. Since (by definition) every element of  $X$  has exactly one arrow, this implies that

$$|X| = \sum_{y \in Y} d(y).$$

For part (a), suppose that  $f : X \rightarrow Y$  is surjective, i.e., that we have  $d(y) \geq 1$  for all  $y \in Y$ . In this case we have

$$|X| = \sum_{y \in Y} d(y) \geq \sum_{y \in Y} 1 = |Y|.$$

For part (b), suppose that  $f : X \rightarrow Y$  is injective, i.e., that we have  $d(y) \leq 1$  for all  $y \in Y$ . In this case we have

$$|X| = \sum_{y \in Y} d(y) \leq \sum_{y \in Y} 1 = |Y|.$$

For part (c), suppose that  $f : X \rightarrow Y$  is bijective, i.e., that we have  $d(y) = 1$  for all  $y \in Y$ . In this case we have

$$|X| = \sum_{y \in Y} d(y) = \sum_{y \in Y} 1 = |Y|.$$

□

**Problem 2.** For all integers  $a, b \in \mathbb{Z}$  with  $b \neq 0$ , we define an **abstract symbol** “ $\frac{a}{b}$ ”. We declare rules for “multiplying” and “adding” abstract symbols,

$$\frac{a}{b} \cdot \frac{c}{d} := \frac{ac}{bd} \quad \text{and} \quad \frac{a}{b} + \frac{c}{d} := \frac{ad + bc}{bd},$$

and we declare that the abstract symbols  $\frac{a}{b}$  and  $\frac{c}{d}$  are “equal” if and only if  $ad = bc$ . Let  $\mathbb{Q}$  denote the set of abstract symbols (we call this the system of **rational numbers**). For all rational numbers  $x \in \mathbb{Q}$ , prove that  $x$  can be expressed as  $\frac{a}{b}$  where  $a, b \in \mathbb{Z}$  have no common divisor except  $\pm 1$ . (We say that the fraction  $x$  can be written in “lowest terms”.) [Hint: Let  $S$  be the set of absolute values of all the possible numerators of  $x$ :

$$S := \left\{ |a| \in \mathbb{N} : \exists a, b \in \mathbb{Z} \text{ such that } x = \frac{a}{b} \right\} \subseteq \mathbb{N}.$$

Since  $x \in \mathbb{Q}$ , the set  $S$  is not empty, so by Well-Ordering it has a smallest element.]

*Proof.* Consider a rational number  $x \in \mathbb{Q}$ , and let  $S$  be the set of absolute values of all possible numerators of  $x$ . That is, let

$$S := \left\{ |a| : \exists a, b \in \mathbb{Z} \text{ such that } x = \frac{a}{b} \right\}$$

Note that  $S$  is a subset of the natural numbers  $\mathbb{N}$ . Since  $x \in \mathbb{Q}$ , we know that  $x$  can be expressed as a fraction in at least one way, hence  $S \neq \emptyset$ . Thus, by the Well-Ordering Principle  $S$  has a smallest element. Call it  $m \in S$ .

Note briefly that for all  $a, b \in \mathbb{Z}$  with  $b \neq 0$  we have  $\frac{-a}{b} = \frac{a}{-b}$ . Thus the possible numerators of  $x$  come in positive-negative pairs. Since  $m \in S$  we conclude that there exists  $n \in \mathbb{Z}$  such that  $x = \frac{m}{n}$ . Now we claim that  $m$  and  $n$  have no nontrivial common divisor, i.e., that  $\gcd(m, n) = 1$ . To prove this, assume for contradiction that there exists  $d \in \mathbb{Z}$  such that  $d|m$  (say  $m = dm'$ ),  $d|n$  (say  $n = dn'$ ), and  $|d| > 1$ . Then we have

$$x = \frac{m}{n} = \frac{dm'}{dn'} = \frac{m'}{n'}$$

and we see that  $m'$  is also an element of  $S$ . But since  $|d| > 1$  we have  $|m| = |dm'| = |d||m'| > |m'|$ , which contradicts the minimality of  $m$ . We conclude that  $m, n$  have no common divisor, thus we have succeeded in writing  $x$  in lowest terms.  $\square$

*Alternative Proof.* Consider a rational number  $x \in \mathbb{Q}$ . By definition this means that  $x = \frac{a}{b}$  for some  $a, b \in \mathbb{Z}$  with  $b \neq 0$ . Let  $d = \gcd(a, b)$ , with  $a = da'$  and  $b = db'$ , so we can write

$$x = \frac{a}{b} = \frac{da'}{db'} = \frac{a'}{b'}$$

We claim that  $\gcd(a', b') = 1$ . Indeed, by Bézout's Identity there exist  $x, y \in \mathbb{Z}$  such that  $d = ax + by$  and then we have

$$\begin{aligned} d &= ax + by \\ d &= da'x + db'y \\ d &= d(a'x + b'y) \\ 1 &= a'x + b'y. \end{aligned}$$

This means that any common divisor of  $a'$  and  $b'$  also divides 1, hence  $\gcd(a', b') = 1$ . We have thus succeeded in expressing  $x$  in lowest terms.  $\square$

**Problem 3.** The Division Algorithm 2.12 says that for all  $a, b \in \mathbb{Z}$  with  $b > 0$  there exist unique  $q, r \in \mathbb{Z}$  such that  $a = qb + r$  and  $0 \leq r < b$ . Explicitly use this to prove the following: For all  $a, b \in \mathbb{Z}$  with  $b > 0$  there exists a unique integer  $k \in \mathbb{Z}$  such that

$$k \leq \frac{a}{b} < k + 1.$$

[Note: You must prove both the *existence* and the *uniqueness* of  $k$ . Don't be a hero; **quote** the Division Algorithm. You do not need to reduce everything to the axioms.]

*Proof.* First we will prove that existence of  $k \in \mathbb{Z}$ . Applying the Division Algorithm to divide  $a$  by  $b$  yields  $a = qb + r$  with  $0 \leq r < b$ . Then we have

$$\begin{aligned} 0 &\leq r < b \\ 0 &\leq a - qb < b \\ qb &\leq a < b + qb \\ q &\leq \frac{a}{b} < q + 1. \end{aligned}$$

We may now take  $k = q$ .

Next we will show that this  $k$  is unique. That is, suppose that we have  $k_1 \leq \frac{a}{b} < k_1 + 1$  and  $k_2 \leq \frac{a}{b} < k_2 + 1$ . We want to show that  $k_1 = k_2$ . By reversing the steps above we have

$$\begin{aligned} k_1 &\leq \frac{a}{b} < k_1 + 1 \\ k_1 b &\leq a < k_1 b + b \\ 0 &\leq a - k_1 b < b. \end{aligned}$$

If we let  $r_1 := a - k_1 b$  then we have  $a = k_1 b + r_1$  with  $0 \leq r_1 < b$ . Similarly, if we let  $r_2 := a - k_2 b$  then we have  $a = k_2 b + r_2$  with  $0 \leq r_2 < b$ . By the uniqueness part of the Division Algorithm this implies that  $k_1 = k_2$ , as desired.  $\square$

**Problem 4. How do  $-$  and  $\times$  interact?** Prove the following exercises using the axioms of  $\mathbb{Z}$  from the handout. It will save time if you assume the Cancellation Property that was proved on the previous homework:  $\forall a, b, c \in \mathbb{Z}, (a + b = a + c) \Rightarrow (b = c)$ .

- (a) Prove that for all  $a \in \mathbb{Z}$  we have  $0a = 0$ .
- (b) Recall that  $-n$  is the unique integer such that  $n + (-n) = 0$ . Prove that for all  $a, b \in \mathbb{Z}$  we have  $(-a)b = -(ab)$ . [Hint: You will need part (a).]
- (c) Prove that for all  $a, b, c \in \mathbb{Z}$  we have  $a(b - c) = ab - ac$ . [Hint: Use part (b).]
- (d) Prove that for all  $a, b \in \mathbb{Z}$  we have  $(-a)(-b) = ab$ . [Hint: Use part (a) to show that  $ab + a(-b) = 0$  and then use part (b). Note that  $-(-n) = n$  for all  $n \in \mathbb{Z}$ .]

[Now if a child asks you **why** negative times negative is positive, you will know what to say.]

*Proof.* I will apply the commutative axioms (A1) and (M1) when needed, without comment. To prove (a) first note that  $0 = 0 + 0$  by axiom (A3). Then we have

$$\begin{aligned} 0 &= 0 + 0, \\ 0a &= (0 + 0)a, \\ 0a &= 0a + 0a, & \text{(D)} \\ 0 + 0a &= 0a + 0a. & \text{(A3)} \end{aligned}$$

Cancelling  $0a$  from the final equation gives  $0 = 0a$ . To prove (b), recall that  $-(ab)$  is the unique integer  $x$  such that  $ab + x = 0$ . Thus we need to show that  $ab + (-a)b = 0$ . Indeed, we have

$$\begin{aligned} ab + (-a)b &= (a + (-a))b, & \text{(D)} \\ &= 0b, & \text{(A3)} \\ &= 0. & \text{by part (a)} \end{aligned}$$

To prove (c) note that

$$\begin{aligned} a(b - c) &= a(b + (-c)), \\ &= ab + a(-c), && \text{(D)} \\ &= ab + (-(ac)), && \text{by part (b)} \\ &= ab - ac. \end{aligned}$$

Finally, to prove (d) first note that

$$\begin{aligned} ab + a(-b) &= a(b + (-b)), && \text{(D)} \\ &= a0, && \text{(A3)} \\ &= 0. && \text{by part (a)} \end{aligned}$$

This means that  $ab$  is the additive inverse of  $a(-b)$ , i.e.  $ab = -(a(-b))$ . Recall that the result of part (b) says that  $-(xy) = (-x)y$  for all  $x, y \in \mathbb{Z}$ . We apply this with  $x = a$  and  $y = -b$  to conclude that

$$ab = -(a(-b)) = (-a)(-b).$$

□