

3.1. Let \mathcal{I} be a small category and let \mathcal{C} be any category. For each object $c \in \mathcal{C}$ we define the *constant diagram* $\Delta(c) : \mathcal{I} \rightarrow \mathcal{C}$ which sends each object of \mathcal{I} to c and each arrow of \mathcal{I} to id_c .

- (a) Verify that the map $c \mapsto \Delta(c)$ defines a functor $\Delta : \mathcal{C} \rightarrow \mathcal{C}^{\mathcal{I}}$, called the *diagonal functor*.
- (b) Consider any diagram $D : \mathcal{I} \rightarrow \mathcal{C}$ and object $c \in \mathcal{C}$. Verify that a natural transformation $\alpha : \Delta(c) \Rightarrow D$ is the same thing we previously called a “cone under D ”.

3.2. State the definition for the *Cartesian product* of categories. For any category \mathcal{C} , verify that the assignment $(x, y) \mapsto \text{Hom}_{\mathcal{C}}(x, y)$ defines a functor

$$\text{Hom}_{\mathcal{C}}(-, -) : \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \text{Set}.$$

3.3. Verify that the ugly and fancy definitions of adjoint functors in the notes are equivalent.

3.4. Work through the details of the Yoneda Lemma with me.