

4/29/16

HW 4 due now.

Final Exam next Wed May 4,
11:00-1:30 in this room (Ungar 411).

Today: Review for Final.

The Final Exam will cover the material discussed in lecture since the Midterm, as well as the material from HW 3 & HW 4. Here are the main topics.

① The Category $R\text{-Alg}$.

Let R be a commutative ring. We define an R -algebra as a pair (S, φ) where

- S is a (possibly noncommutative) ring.
- $\varphi: R \rightarrow S$ is a ring homomorphism
- in $\varphi \in Z(S) = \{s \in S : st = ts \forall t \in S\}$

We define a morphism of R -algebras

$\Phi: (S_1, \varphi_1) \rightarrow (S_2, \varphi_2)$ as a ring homomorphism $\Phi: S_1 \rightarrow S_2$ such that the following triangle commutes:

$$\begin{array}{ccc}
 S_1 & \xrightarrow{\Phi} & S_2 \\
 \nwarrow \varphi_1 & & \nearrow \varphi_2 \\
 & R &
 \end{array}$$

One can check that this defines a category, which we call $R\text{-Alg}$. Note that we have a forgetful functor

$$U: R\text{-Alg} \rightarrow R\text{-Mod}$$

defined by "forgetting the monoid structure". Specifically, if $\varphi: R \rightarrow S$ is an R -algebra then we obtain a ring homomorphism

$$\lambda: R \rightarrow \text{End}_{\text{Ab}}(|S|)$$

by sending $r \in R$ to the "multiplication map" $\lambda_r(s) := \varphi(r)s = s\varphi(r)$. One can check that this defines an R -module structure on the abelian group $|S|$. [In fact, you did check this on HW 2.5(b).] With respect to this structure we can define a subring

$$\text{End}_R(|S|) \subseteq \text{End}_{\text{Ab}}(|S|),$$

and then observe ("trivially" but not necessarily "easily") that the ring homomorphism λ sends R into the center of this subring:

$$\lambda: R \rightarrow \text{im } \lambda \subseteq Z(\text{End}_R(|S|)).$$

[We did this in class on 4/7/16].

Thus we obtain an R -algebra structure on the ring $\text{End}_R(|S|)$ that is compatible with the R -algebra structure on the ring S in the sense that for all $s_1, s_2, s_3 \in S$ and $r \in R$ we have

- $s_1(s_2 + \lambda_r(s_3)) = s_1 s_2 + \lambda_r(s_1 s_3)$
- $(s_1 + \lambda_r(s_2)) s_3 = s_1 s_3 + \lambda_r(s_2 s_3)$.

[We say that the monoid structure on S is " R -bilinear".]

[Remark: The forgetful functor

$U: R\text{-Alg} \rightarrow R\text{-Mod}$ has a left-adjoint free functor $T: R\text{-Mod} \rightarrow R\text{-Alg}$ called the "tensor algebra" $M \mapsto T(M)$.



Unfortunately, we didn't have time to get into this.]

More generally, if M is any R -module defined by a ring homomorphism

$$\lambda: R \rightarrow \text{End}_{\text{Ab}}(M)$$

then since R is commutative we automatically ("tautologically") obtain an R -algebra of R -linear endomorphisms

$$\lambda: R \rightarrow \text{im } \lambda \subseteq Z(\text{End}_R(M)).$$

In this sense, endomorphisms of R -modules are the "concrete" prototype for R -algebras, and thus tell us how we should define "representations" of algebras.

This also motivates the following, more "concrete", definition of R -algebras.



★ Concrete Definition:

Let R be a commutative ring. An R -algebra is a structure (M, λ, \circ, id) where

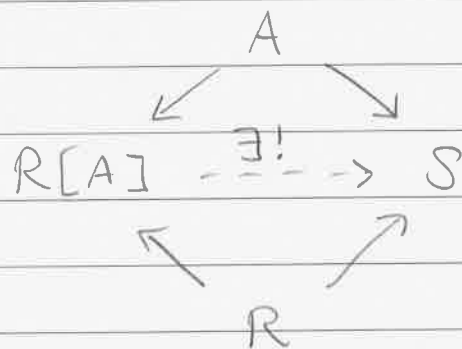
- (M, λ) is an R -module
- (M, \circ, id) is a monoid structure making M into a ring (i.e. \circ distributes over addition in M)
- The operation $\circ: M \times M \rightarrow M$ is compatible with λ (i.e. it is R -bilinear).

[Any book on representation theory expects you to know this, but no algebra text that I know of actually teaches this.]

② The category $R\text{-CAlg}$.

If we restrict our attention to R -algebras $\varphi: R \rightarrow S$ in which S is a commutative ring then we obtain a full subcategory $R\text{-CAlg} \subseteq R\text{-Alg}$, called the category of commutative R -algebras. This was the subject of HW3.

Recall that for any set $A \in \text{Set}$ there exists a free commutative algebra $A \rightarrow R[A]$ satisfying the universal property



[There are too many maps here to give them names, so we'll have to be careful!]

If $|A| = n$ then we can identify $R[A]$ with the polynomial algebra in n commuting variables, called

$$R[x_1, x_2, \dots, x_n]$$

[Formally we can view this as the "symmetric algebra" $S(R^{\oplus A})$ of the free module $R^{\oplus A}$. This $S: R\text{-Mod} \rightarrow R\text{-Alg}$ is the left adjoint to the forgetful functor $U: R\text{-Alg} \rightarrow R\text{-Mod}$. If $F: \text{Set} \rightarrow R\text{-Mod}$ is the free module functor then we obtain

$$R[A] = S(F(A)).$$



Unfortunately, we didn't have time to get into this.]

Now suppose that $A = \{s_1, s_2, \dots, s_n\} \xrightarrow{i} S$.

Then we obtain the canonical "evaluation morphism" from the free algebra

$$\begin{aligned}\varphi_A : R[x_1, \dots, x_n] &\rightarrow S \\ f(x_1, \dots, x_n) &\mapsto f(s_1, \dots, s_n).\end{aligned}$$

We say that

- $A \hookrightarrow S$ is R -algebraically independent if φ_A is injective
- $A \hookrightarrow S$ is R -algebraically generating if φ_A is surjective.

We also sometimes use the notation

$$R[A] = \text{im } \varphi_A \subseteq S$$

but be warned: this notation conflicts with the previous usage because $\text{im } \varphi_A$ is a free algebra only when $\ker \varphi_A = 0$. We'll try not to get confused.

Furthermore, we say that

- $A \hookrightarrow S$ is R -algebraically almost-generating if S is algebraic over the subring $R[A] = \text{im } \varphi_A$.
[This means that for all $s \in S$ the evaluation morphism $\varphi_s: R[A][x] \rightarrow S$ is non-injective.]
- $A \hookrightarrow S$ is an R -transcendence basis if it is algebraically independent and almost-generating.

★ Steinitz Exchange for Algebras :

Let $S \in R\text{-CAlg}$ be an integral domain, let $A \subseteq S$ be algebraically independent and let $B \subseteq S$ be almost-generating. Then we can replace elements of B by elements of A while maintaining the almost-generating property. It follows that

$$|A| \leq |B|$$

and hence all R -transcendence bases have the same cardinality, called the R -transcendence degree of S .

No, you do not need to memorize the proof of this theorem, but you should understand the proof [in the HW3 solutions] and be able to reproduce bits and pieces of it.

Here are the key lemmas.

Lemma 1: If $K_1 \subseteq K_2 \subseteq K_3$ are fields then

$$[K_3:K_1] = [K_3:K_2][K_2:K_1].$$

Lemma 2: Let $K \subseteq L$ be fields.

- If $[L:K] < \infty$ then $K \subseteq L$ is algebraic.
- If $K \subseteq L$ is algebraic and finitely generated as a K -algebra then $[L:K] < \infty$.

Lemma 3: Let $K \subseteq L$ be fields. Then the algebraic closure $K \subseteq \text{Alg}_K(L) \subseteq L$ is a subring of L (in fact a subfield).

Lemma 4: Let $K_1 \subseteq K_2 \subseteq K_3$ be fields. If $\text{Alg}_{K_1}(K_2) = K_2$ and $\text{Alg}_{K_2}(K_3) = K_3$, then $\text{Alg}_{K_1}(K_3) = K_3$.

}

Lemma 5: Lemmas 3 & 4 also hold for integral domains by the trick of "localization".

[This whole story generalizes to situations where R is an integral domain but S need not be (see "Noether Normalization"). However, the proofs and statements become more difficult.]

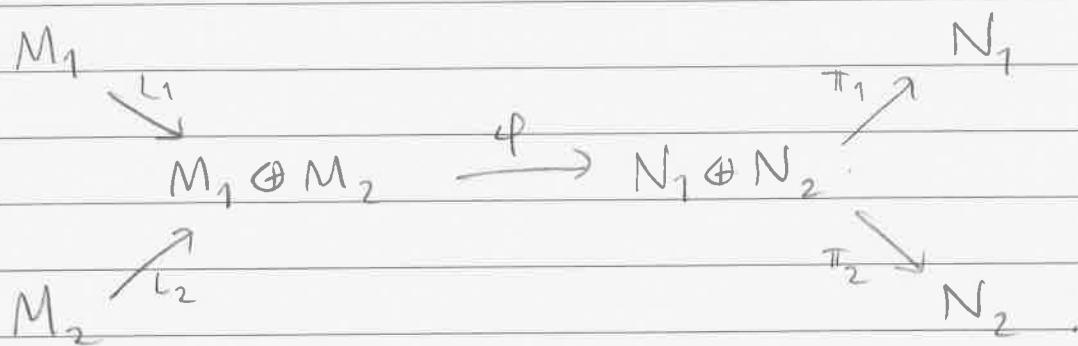
(3) Matrix Notation.

Recall that finite products and coproducts coincide in the category $R\text{-Mod}$ (more generally, in any "additive category"). We denote this "biproduct" by \oplus .

Now consider any homomorphism between direct sums of modules

$$\varphi: M_1 \oplus M_2 \longrightarrow N_1 \oplus N_2.$$

By interpreting $N_1 \oplus N_2$ as a product and $M_1 \oplus M_2$ as a coproduct we obtain a diagram:



By the universal properties of product & coproduct, the map φ is uniquely determined by its four "components"

$$\varphi_{ij} := \pi_i \circ \varphi \circ L_j : M_j \rightarrow N_i.$$

Furthermore, we can express φ in terms of the following matrix notation. Given elements $m = m_1 + m_2 \in M_1 \oplus M_2$ and $n = n_1 + n_2 \in N_1 \oplus N_2$ we will write

$$m = \begin{pmatrix} m_1 \\ m_2 \end{pmatrix} \quad \& \quad n = \begin{pmatrix} n_1 \\ n_2 \end{pmatrix}$$

and then we have

$$\varphi(m) = \begin{pmatrix} \varphi_{11} & \varphi_{12} \\ \varphi_{21} & \varphi_{22} \end{pmatrix} \begin{pmatrix} m_1 \\ m_2 \end{pmatrix} = \begin{pmatrix} \varphi_{11}(m_1) + \varphi_{12}(m_2) \\ \varphi_{21}(m_1) + \varphi_{22}(m_2) \end{pmatrix}.$$

Now suppose we have a homomorphism between free modules

$$\varphi : R^{\oplus A} \rightarrow R^{\oplus B}$$

If the bases are $A = \{a_1, \dots, a_m\}$ and $B = \{b_1, \dots, b_n\}$ then we can write this as

$$\varphi : \bigoplus_{j=1}^m (a_j) \rightarrow \bigoplus_{i=1}^n (b_i)$$

and we can define the component morphisms

$$\varphi_{ij} := \pi_i \circ \varphi \circ \iota_j : (a_j) \rightarrow (b_i)$$

If we denote an element $r = \sum_j r_j a_j \in R^{\oplus A}$ as a column vector

$$r = \begin{pmatrix} r_1 \\ r_2 \\ \vdots \\ r_m \end{pmatrix}$$

then by iterating the construction above we can express φ in matrix notation as follows :

$$\varphi(r) = \begin{pmatrix} \varphi_{11} & \dots & \varphi_{1m} \\ \vdots & & \vdots \\ \varphi_{n1} & \dots & \varphi_{nm} \end{pmatrix} \begin{pmatrix} r_1 \\ \vdots \\ r_m \end{pmatrix} = \begin{pmatrix} \sum_j \varphi_{1j}(r_j) \\ \vdots \\ \sum_j \varphi_{nj}(r_j) \end{pmatrix}.$$

Since each component $\varphi_{ij} : (a_j) \rightarrow (b_i)$ is a homomorphism of rank 1 free modules it is uniquely determined by the ring element $s_{ij} \in R$ such that

$$\varphi_{ij}(a_j) = s_{ij} b_i.$$

Thus it is more meaningful to write

$$\varphi = \begin{pmatrix} s_{11} & \dots & s_{1m} \\ \vdots & & \vdots \\ s_{n1} & \dots & s_{nm} \end{pmatrix}.$$

By composing homomorphisms between free modules we recover the familiar definition of matrix multiplication. In summary, if F_1, F_2 are finitely generated free R -modules then any choice of bases

$$F_1 = R^{\oplus A} \quad \& \quad F_2 = R^{\oplus B}$$

}

determines an isomorphism of abelian groups

$$\Phi_{AB} : \text{Hom}_R(F_1, F_2) \xrightarrow{\sim} \text{Mat}_{|B| \times |A|}(R)$$

an isomorphism of rings

$$\Phi_A : \text{End}_R(F_1) \xrightarrow{\sim} \text{Mat}_{|A|}(R)$$

and an isomorphism of groups

$$\Phi_A : \text{Aut}_R(F_1) \xrightarrow{\sim} \text{GL}_{|A|}(R).$$

[If, moreover, R is commutative then the first is an isomorphism of R -modules & the second is an isomorphism of R -algebras.]

The great benefit of matrix notation is that we can use Gaussian elimination to obtain constructive proofs and a more concrete understanding of modules (at least over nice rings).

Specifically, we define the elementary matrices $E_{ij}(r)$, $E_{ii}(r)$, $P_{ij} \in \text{GL}_n(R)$. If R is a Euclidean domain then these generate the group $\text{GL}_n(R)$.

If R is, more generally, a PID then we also need the pseudo-elementary matrices:

$$E_{ij} \begin{pmatrix} a & b \\ c & d \end{pmatrix} := I + (a-1)e_{ii} + be_{ij} + ce_{ji} + (d-1)e_{jj}.$$

Now we can describe the action of the group $GL_m(R) \times GL_n(R)$ on the set $\text{Mat}_{m \times n}(R)$ by left and right multiplication

$$(P, Q) \cdot A := PAQ^{-1}$$

in terms of algorithms involving elementary (and pseudo-elementary) row and column operations. This leads to the following theorem.

★ Smith Normal Form:

Let R be a PID and consider a matrix $A \in \text{Mat}_{m \times n}(R)$. Then there exist invertible matrices

$$P \in GL_m(R) \quad \& \quad Q \in GL_n(R)$$

and ideals $(r_d) \subseteq \dots \subseteq (r_2) \subseteq (r_1) \subseteq R$ such that

$$PAQ^{-1} = \left(\begin{array}{ccc|c} r_1 & & & 0 \\ & r_2 & & \\ & & \ddots & \\ & & & r_d \\ \hline & & & 0 \\ & 0 & & 0 \end{array} \right)$$

Furthermore, the matrices P, Q and elements r_1, r_2, \dots, r_d are unique up to multiplication by units. We call PAQ^{-1} the Smith Normal Form of A .

(4) FTFGMPID, Part II.

This allows us to completely determine the structure of finitely generated modules over a PID. Let R be a PID and let $M \in R\text{-Mod}$ be finitely generated. Then by results from before the Spring Break we obtain a short exact sequence

$$0 \rightarrow R^{\oplus m} \xrightarrow{\varphi} R^{\oplus n} \rightarrow M \rightarrow 0,$$

so that $M \cong R^{\oplus n} / \text{im } \varphi$.

⑤ $K[x]$ -modules.

The FTFGMPID has an immediate application to abelian groups (\mathbb{Z} -modules) but I think the more interesting applications relate to the "other" PID (namely, $K[x]$ where K is a field).

Using the fact that $K[x]$ is the free K -algebra generated by one element gives us a bijection between $K[x]$ -modules and pairs (V, φ) where

- V is a K -vector space
- $\varphi \in \text{End}_K(V)$ is a K -linear endomorphism

If V is finite dimensional (say $V = K^{\oplus n}$) then by choosing a basis we can think of φ as an $n \times n$ matrix $A = [\varphi] \in \text{Mat}_n(K)$. Furthermore, two such modules (V_1, A_1) & (V_2, A_2) are isomorphic as $K[x]$ -modules if and only if there exists an invertible matrix $P \in \text{GL}_n(K)$ such that

$$PA_1P^{-1} = A_2.$$


Finally, since $K[x]$ is a PID we obtain from the FTFGMPID an isomorphism of $K[x]$ -modules

$$V \cong \bigoplus_{i=1}^d K[x]/(f_i(x))$$

for some unique monic nonconstant polynomials

$$f_1(x) \mid f_2(x) \mid \dots \mid f_d(x).$$

[Recall: $f_d(x)$ is the minimal polynomial & $\prod f_i(x)$ is the characteristic polynomial of the corresponding morphism.]

By choosing the most obvious bases for the direct summands we can then obtain the Rational Canonical Form and the Jordan Canonical Form (when K is algebraically closed) of the corresponding matrix. 

There is more to say but it's been a long semester so I think I'll stop here.