

**Problem 1.** De Morgan's Laws say that for all statements  $P, Q$  we have

$$\neg(P \vee Q) = \neg P \wedge \neg Q \quad \text{and} \quad \neg(P \wedge Q) = \neg P \vee \neg Q.$$

- (a) Use truth tables to prove these laws.
- (b) Use a truth table to prove that  $(P \Rightarrow Q) = (\neg P) \vee Q$  for all statements  $P, Q$ .
- (c) Combine parts (a) and (b) to prove that for all statements  $P, Q$  we have

$$(P \Rightarrow Q) = (\neg Q \Rightarrow \neg P).$$

Do **not** use a truth table.

(a) Here is a truth table proving  $\neg(P \vee Q) = \neg P \wedge \neg Q$ :

$P$	$Q$	$P \vee Q$	$\neg(P \vee Q)$	$\neg P$	$\neg Q$	$\neg P \wedge \neg Q$
$T$	$T$	$T$	$F$	$F$	$F$	$F$
$T$	$F$	$T$	$F$	$F$	$T$	$F$
$F$	$T$	$T$	$F$	$T$	$F$	$F$
$F$	$F$	$F$	$T$	$T$	$T$	$T$

And here is a truth table proving  $\neg(P \wedge Q) = \neg P \vee \neg Q$ :

$P$	$Q$	$P \wedge Q$	$\neg(P \wedge Q)$	$\neg P$	$\neg Q$	$\neg P \vee \neg Q$
$T$	$T$	$T$	$F$	$F$	$F$	$F$
$T$	$F$	$F$	$T$	$F$	$T$	$T$
$F$	$T$	$F$	$T$	$T$	$F$	$T$
$F$	$F$	$F$	$T$	$T$	$T$	$T$

(b) And here is a truth table proving that  $(P \Rightarrow Q) = (\neg P) \vee Q$ :

$P$	$Q$	$P \Rightarrow Q$	$\neg P$	$\neg P \vee Q$
$T$	$T$	$T$	$F$	$T$
$T$	$F$	$F$	$F$	$F$
$F$	$T$	$T$	$T$	$T$
$F$	$F$	$T$	$T$	$T$

(c) *Proof.* From part (b) we know that  $(A \Rightarrow B) = \neg A \vee B$  for all statements  $A$  and  $B$ . Substituting  $A = P$  and  $B = Q$  gives

$$(P \Rightarrow Q) = \neg P \vee Q,$$

and substituting  $A = \neg Q$  and  $B = \neg P$  gives

$$(\neg Q \Rightarrow \neg P) = \neg(\neg Q) \vee (\neg P) = Q \vee \neg P.$$

Since  $\neg P \vee Q = Q \vee \neg P$  we conclude that  $(P \Rightarrow Q) = (\neg Q \Rightarrow \neg P)$ . □

**Problem 2.** Practice with logical analysis.

- (a) Use the results of Problem 1 to prove that for all statements  $P, Q, R$  we have

$$P \Rightarrow (Q \vee R) = (\neg Q \wedge \neg R) \Rightarrow \neg P.$$

Do **not** use a truth table.

(b) Use the result of (a) to prove that for all  $a, b, c, d \in \mathbb{Z}$  we have

$$a + b \leq c + d \implies a \leq c \text{ or } b \leq d.$$

(c) Prove that the converse of the statement in part (b) is **false**. [Hint: To prove that a universal statement is false it is enough to provide a single counterexample.]

(a) *Proof.* For all statements  $P, Q, R$  we have

$$\begin{aligned} P \implies (Q \vee R) &= \neg(Q \vee R) \implies \neg P && \text{by 1(c)} \\ &= (\neg Q \wedge \neg R) \implies \neg P. && \text{by 1(a)} \end{aligned}$$

□

(b) Let  $a, b, c, d \in \mathbb{Z}$  and define the statements

$$P = \text{“} a + b \leq c + d, \text{”}$$

$$Q = \text{“} a \leq c, \text{”}$$

$$R = \text{“} b \leq d. \text{”}$$

I claim that  $P \implies (Q \vee R)$ . *Proof.* By part (a) it is enough to prove that  $(\neg Q \wedge \neg R) \implies \neg P$ . In other words, we will prove that

$$(a > c \text{ and } b > d) \implies (a + b > c + d).$$

So let us assume that  $a > c$  and  $b > d$ . By adding  $b$  to both sides of  $a > c$  we obtain

$$a > c$$

$$a + b > b + c,$$

and by adding  $c$  to both sides of  $b > d$  we obtain

$$b > d$$

$$b + c > c + d.$$

Then by the “transitivity” of the relation “ $>$ ” we conclude that  $a + b > c + d$ . □

(c) The statement from (b) says that

$$\forall a, b, c, d \in \mathbb{Z}, (a + b \leq c + d) \implies (a \leq c \vee b \leq d).$$

and the **converse** of this statement says that

$$\forall a, b, c, d \in \mathbb{Z}, (a \leq c \vee b \leq d) \implies (a + b \leq c + d).$$

To prove that the converse is **false** we will prove that the **opposite** of the converse is true. By applying de Morgan’s law and the property  $(P \not\Rightarrow Q) = \neg(P \implies Q) = P \wedge \neg Q$ , the statement we want to prove is

$$\begin{aligned} &\neg \left[ \forall a, b, c, d \in \mathbb{Z}, (a \leq c \vee b \leq d) \implies (a + b \leq c + d) \right] \\ &= \exists a, b, c, d \in \mathbb{Z}, (a \leq c \vee b \leq d) \not\Rightarrow (a + b \leq c + d) \\ &= \exists a, b, c, d \in \mathbb{Z}, (a \leq c \vee b \leq d) \wedge \neg(a + b \leq c + d) \\ &= \exists a, b, c, d \in \mathbb{Z}, (a \leq c \vee b \leq d) \wedge (a + b > c + d). \end{aligned}$$

In order to prove that such integers exist, it is enough to give one example. So let us choose  $(a, b, c, d) = (1, 2, 1, 1)$ . Then  $(1 \leq 1 \vee 2 \leq 1)$  is a true statement because at least one of the statements  $1 \leq 1$  and  $2 \leq 1$  is true, and the statement  $(1 + 2 > 1 + 1)$  is also true. □

[Remark: It is okay for you to skip some of the logical analysis and jump right to the counterexample, but I wanted to show all the gory details for pedagogical reasons.]

**Problem 3.** I will guide you through an induction proof that

$$(a - 1) \mid (a^n - 1) \quad \text{for all integers } a, n \in \mathbb{Z} \text{ such that } n \geq 1.$$

For the purpose of the proof, let  $a \in \mathbb{Z}$  be a fixed integer. We will use induction on  $n$ .

- (a) Prove that  $(a - 1) \mid (a^n - 1)$  when  $n = 1$ .
- (b) Now assume that  $(a - 1) \mid (a^n - 1)$  is true for some fixed  $n \geq 1$ . In this case, prove that

$$(a - 1) \mid (a^{n+1} - 1).$$

(a) *Proof.* Since  $(a - 1) = (a - 1) \cdot 1$  and  $1 \in \mathbb{Z}$  we conclude that  $(a - 1) \mid (a - 1)$ . □

(b) Fix a positive integer  $n \geq 1$  and **assume** that  $(a - 1) \mid (a^n - 1)$ . By definition this means that there exists some  $k \in \mathbb{Z}$  such that  $(a^n - 1) = (a - 1)k$ . Now multiply both sides of this equation by  $a$  to get

$$\begin{aligned} (a^n - 1) &= (a - 1)k \\ (a^{n+1} - a) &= (a - 1)ak, \end{aligned}$$

and then add  $(a - 1)$  to both sides to get

$$\begin{aligned} (a^{n+1} - a) &= (a - 1)ak \\ (a^{n+1} - a) + (a - 1) &= (a - 1)ak + (a - 1) \\ (a^{n+1} - 1) &= (a - 1)(ak + 1). \end{aligned}$$

Since  $(ak + 1) \in \mathbb{Z}$  we conclude that  $(a - 1) \mid (a^{n+1} - 1)$  as desired. □

[Remark: It follows, for example, that  $58^{100} - 1$  is a multiple of 57.]

**Problem 4.** For all integers  $d \in \mathbb{Z}$  let us define the statement

$$P(d) := “ \forall n \in \mathbb{Z}, d \mid n^2 \Rightarrow d \mid n. ”$$

- (a) Now fix an integer  $d \geq 2$  and prove that

$$P(d) \implies \sqrt{d} \notin \mathbb{Q}$$

[Hint: Mimic the proofs from class when  $d = 2$  and  $d = 3$ .]

- (b) Prove that  $P(5)$  is a true statement, and hence that  $\sqrt{5}$  is irrational.
- (c) Prove that  $P(12)$  is a false statement. [Remark: It is still true that  $\sqrt{12}$  is irrational, but the method of proof from part (a) will not work. Maybe you can see how to fix it.]

(a) *Proof that  $P(d) \implies \sqrt{d} \notin \mathbb{Q}$ .* Fix an integer  $d \geq 2$  and let us assume that  $P(d)$  is a true statement. In other words, let us assume that

$$(*) \quad d \mid n^2 \implies d \mid n \quad \text{for all } n \in \mathbb{Z}.$$

In this case we will prove that  $\sqrt{d} \notin \mathbb{Q}$ . So let us **assume for contradiction** that  $\sqrt{d} \in \mathbb{Q}$ . Then we can write  $\sqrt{d} = a/b$  where  $a$  and  $b$  are integers with no common factors other than  $\pm 1$ . Now square both sides to get

$$\begin{aligned}d &= a^2/b^2 \\ db^2 &= a^2.\end{aligned}$$

Since  $b^2 \in \mathbb{Z}$  this last equation implies that  $d|b^2$  and then from (\*) we get  $d|n$ . In other words, we have  $n = dk$  for some  $k \in \mathbb{Z}$ . Substituting into the previous equation gives

$$\begin{aligned}db^2 &= a^2 \\ db^2 &= (dk)^2 \\ db^2 &= d^2k^2 \\ b^2 &= da^2.\end{aligned}$$

Since  $a^2 \in \mathbb{Z}$  this last equation tells us that  $d|b^2$  and then from (\*) we get  $d|b$ . In summary, we have shown that  $d|a$  and  $d|b$ , which contradicts the fact that  $a$  and  $b$  have no common factor. We conclude that  $\sqrt{d} \in \mathbb{Q}$  is false, and hence  $\sqrt{d} \notin \mathbb{Q}$ .  $\square$

But is the statement  $P(d)$  true or false?

(b)  $P(5)$  is true, hence it follows from part (a) that  $\sqrt{5} \notin \mathbb{Q}$ . *Proof.* For all  $n \in \mathbb{Z}$  we want to show that  $5|n^2$  implies  $5|n$  and we will do this by showing the contrapositive statement:

$$5 \nmid n \Rightarrow 5 \nmid n^2.$$

So consider any  $n \in \mathbb{Z}$  and assume that  $5 \nmid n$ . There are four cases:

- If  $n = 5k + 1$  for some  $k \in \mathbb{Z}$  then we have

$$n^2 = (5k + 1)^2 = 25k^2 + 10k + 1 = 5(5k^2 + 2k) + 1.$$

- If  $n = 5k + 2$  for some  $k \in \mathbb{Z}$  then we have

$$n^2 = (5k + 2)^2 = 25k^2 + 20k + 4 = 5(5k^2 + 4k) + 4.$$

- If  $n = 5k + 3$  for some  $k \in \mathbb{Z}$  then we have

$$n^2 = (5k + 3)^2 = 25k^2 + 30k + 9 = 5(5k^2 + 6k + 1) + 4.$$

- If  $n = 5k + 4$  for some  $k \in \mathbb{Z}$  then we have

$$n^2 = (5k + 4)^2 = 25k^2 + 40k + 16 = 5(5k^2 + 8k + 3) + 1.$$

In any case, we conclude that  $5 \nmid n^2$ .  $\square$

(c)  $P(12)$  is false. *Proof.* The opposite of the statement  $P(12)$  is

$$\begin{aligned}\neg(\forall n \in \mathbb{Z}, 12|n^2 \Rightarrow 12|n) &= (\exists n \in \mathbb{Z}, 12|n^2 \not\Rightarrow 12|n) \\ &= (\exists n \in \mathbb{Z}, 12|n^2 \text{ but } 12 \nmid n).\end{aligned}$$

To prove the existence of such an integer  $n \in \mathbb{Z}$  it is enough to take  $n = 6$  and observe that

$$12 \mid 36 \quad \text{but} \quad 12 \nmid 6.$$

$\square$

**Bonus Material.**  $\sqrt{12} \notin \mathbb{Q}$ .

*Proof.* We already proved in class that  $\sqrt{3} \notin \mathbb{Q}$  is irrational. Now assume for contradiction that  $\sqrt{12} \in \mathbb{Q}$ . This means we can write  $\sqrt{12} = a/b$  for some integers  $a, b \in \mathbb{Z}$ . But then we have

$$\begin{aligned}\sqrt{12} &= a/b \\ 2\sqrt{3} &= a/b \\ \sqrt{3} &= a/(2b).\end{aligned}$$

Since  $a$  and  $2b$  are integers this implies that  $\sqrt{3} \in \mathbb{Q}$ . Contradiction. □

[Remark: Much later in the course we will see that  $P(d)$  is a true statement for all integers  $d$  that have **no repeated prime factors**. Then by mimicking the proof of the case  $d = 12$  it will follow that

$$\sqrt{d} \notin \mathbb{Z} \Rightarrow \sqrt{d} \notin \mathbb{Q} \quad \text{for all integers } d \in \mathbb{Z}.$$

But you won't have to wait that long because there are easier ways to prove this.]