

There are 4 problems, worth 5 points each. This is a closed book test. Anyone caught cheating will receive a score of **zero**.

**Problem 1.** Let  $P$  and  $Q$  be Boolean variables.

- (a) Draw the truth table for the Boolean function  $P \Rightarrow Q$ .

$P$	$Q$	$P \Rightarrow Q$
$T$	$T$	$T$
$T$	$F$	$F$
$F$	$T$	$T$
$F$	$F$	$T$

- (b) Use a truth table to prove that  $P \Rightarrow Q$  is logically equivalent to the Boolean function  $(\text{NOT } P) \text{ OR } Q$ .

$P$	$Q$	$\text{NOT } P$	$(\text{NOT } P) \text{ OR } Q$	$P \Rightarrow Q$
$T$	$T$	$F$	$T$	$T$
$T$	$F$	$F$	$F$	$F$
$F$	$T$	$T$	$T$	$T$
$F$	$F$	$T$	$T$	$T$

- (c) Apply de Morgan's law and part (b) to find an expression for  $\text{NOT}(P \Rightarrow Q)$ .

$$\begin{aligned} \text{NOT}(P \Rightarrow Q) &= \text{NOT}((\text{NOT } P) \text{ OR } Q) \\ &= (\text{NOT}(\text{NOT } P)) \text{ AND } (\text{NOT } Q) \\ &= P \text{ AND } (\text{NOT } Q). \end{aligned}$$

**Problem 2.** Let  $P$ ,  $Q$  and  $R$  be Boolean variables.

- (a) Use a truth table to prove that the Boolean function  $P \Rightarrow (Q \Rightarrow R)$  is logically equivalent to  $X := (P \text{ AND } (\text{NOT } R)) \Rightarrow (\text{NOT } Q)$ .

$P$	$Q$	$R$	$Q \Rightarrow R$	$P \Rightarrow (Q \Rightarrow R)$	$\text{NOT } R$	$P \text{ AND } (\text{NOT } R)$	$\text{NOT } Q$	$X$
$T$	$T$	$T$	$T$	$T$	$F$	$F$	$F$	$T$
$T$	$T$	$F$	$F$	$F$	$T$	$T$	$F$	$F$
$T$	$F$	$T$	$T$	$T$	$F$	$F$	$T$	$T$
$T$	$F$	$F$	$T$	$T$	$T$	$T$	$T$	$T$
$F$	$T$	$T$	$T$	$T$	$F$	$F$	$F$	$T$
$F$	$T$	$F$	$F$	$T$	$T$	$F$	$F$	$T$
$F$	$F$	$T$	$T$	$T$	$F$	$F$	$T$	$T$
$F$	$F$	$F$	$T$	$T$	$T$	$F$	$T$	$T$

- (b) Now let  $m$  and  $n$  be integers and consider the following statement: “If  $m$  is **odd**, then whenever  $mn$  is **even**, it follows that  $n$  is **even**.” Define  $P$ ,  $Q$  and  $R$  so this statement has the form  $P \Rightarrow (Q \Rightarrow R)$ .

Let  $P = “m$  is odd”,  $Q = “mn$  is even” and  $R = “n$  is even.”

- (c) Use part (a) to prove the statement from part (b).

*Proof.* Instead of the statement  $P \Rightarrow (Q \Rightarrow R)$ , we will prove the statement  $(P \text{ AND } (\text{NOT } R)) \Rightarrow (\text{NOT } Q)$ , which is logically equivalent by part (a). That is, we will prove that for all integers  $m, n \in \mathbb{Z}$  we have  $(m \text{ is odd and } n \text{ is odd}) \Rightarrow (mn \text{ is odd})$ .

So suppose that  $m$  and  $n$  are both odd, say  $m = 2k + 1$  and  $n = 2\ell + 1$  for some integers  $k, \ell \in \mathbb{Z}$ . Then we have

$$\begin{aligned} mn &= (2k + 1)(2\ell + 1) \\ &= 4k\ell + 2k + 2\ell + 1 \\ &= 2(2k\ell + k + \ell) + 1, \end{aligned}$$

which is odd, as desired. □

**Problem 3.** Let  $n$  be an integer and **assume** for the moment the following fact: “If  $n^2$  is a multiple of 5, then so is  $n$ .” Use this fact (i.e. **quote** it at the appropriate time) to **prove by contradiction** that  $\sqrt{5}$  is **not** a ratio of integers.

*Proof.* Suppose for contradiction that  $\sqrt{2}$  is a ratio of integers. Then we can write  $\sqrt{2} = a/b$  where  $a, b \in \mathbb{Z}$  are integers with no common divisor (i.e. the fraction is in “lowest terms”). Squaring both sides gives  $2 = a^2/b^2$  and then multiplying by  $b^2$  gives  $a^2 = 2b^2$ . Now we see that  $a^2$  is even, and the FACT implies that  $a$  is even, say  $a = 2k$ . Substituting this into  $a^2 = 2b^2$  gives  $4k^2 = 2b^2$ , or  $2k^2 = b^2$ . Hence  $b^2$  is even, and the FACT implies that  $b$  itself is even. We now have that  $a$  and  $b$  are **both** even, but this contradicts the assumption that they have no common divisor. □

**Problem 4.** Let  $n$  be an integer.

- (a) What is the contrapositive of this statement?

“If  $n^2$  is a multiple of 5, then so is  $n$ .”

Let  $P = “n^2$  is a multiple of 5” and let  $Q = “n$  is a multiple of 5.” Then the contrapositive of  $P \Rightarrow Q$  is  $(\text{NOT } Q) \Rightarrow (\text{NOT } P)$ , which says: “If  $n$  is **not** a multiple of 5, then **neither** is  $n^2$ .”

- (b) If you want to prove the statement from (a), the argument will break into **four cases**. Tell me what the four cases are.

To prove the statement, we assume that  $n$  is **not** a multiple of 5. There are four ways this can happen: (1)  $n = 5k + 1$  for some  $k \in \mathbb{Z}$ , (2)  $n = 5k + 2$  for some  $k \in \mathbb{Z}$ , (3)  $n = 5k + 3$  for some  $k \in \mathbb{Z}$ , and (4)  $n = 5k + 4$  for some  $k \in \mathbb{Z}$ .

- (c) [**1 Bonus Point**] Finish the proof of the statement from (a). (Use the back of the page if necessary.)

Squaring  $n$  in each case gives: (1)  $n^2 = 5(5k^2 + 2k) + 1$ , (2)  $n^2 = 5(5k^2 + 4k) + 4$ , (3)  $n^2 = 5(5k^2 + 6k + 1) + 4$ , and (4)  $n^2 = 5(5k^2 + 8k + 3) + 1$ . In any case, we conclude that  $n^2$  is **not** a multiple of 5, as desired.