Problem 1. Chinese Remainder Theorem. Find all integers $c \in \mathbb{Z}$ satisfying the following system of simultaneous congruences:

$$\begin{cases} c \equiv 3 \mod 5, \\ c \equiv 6 \mod 9, \\ c \equiv 8 \mod 11 \end{cases}$$

There are two ways to do this.

Two at a time. Recall that the general solution to $c \equiv a \mod m$ and $c \equiv b \mod n$ with gcd(m,n) = 1 is $c \equiv any + bmx \mod mn$, where $x, y \in \mathbb{Z}$ are any integers satisfying mx + ny = 1. First we consider $c \equiv 5 \mod 5$ and $c \equiv 6 \mod 9$. In this case we have (a,b) = (3,6) and (m,n) = (5,9). We observe that the integers (x,y) = (2,-1) satisfy mx + ny = 1. Therefore the general solution is

$$c \equiv any + bmx = 3 \cdot 9 \cdot (-1) + 6 \cdot 5 \cdot 2 \equiv 33 \mod 45.$$

Next we consider the two congruences $c \equiv 8 \mod 11$ and $c \equiv 33 \mod 45$. This time we have (a,b) = (8,33) and (m,n) = (11,45), and we observe that the integers (x,y) = (-4,1) satisfy mx + ny = 1. Therefore the general solution is

$$c \equiv any + bmx = 8 \cdot 45 \cdot 1 + 33 \cdot 11 \cdot (-4) \equiv -1092 \equiv 393 \mod 495.$$

All at once. Alternatively, recall that for any integers satisfying $gcd(m_1, m_2, m_3) = 1$, there exist some integers $x_1, x_2, x_3 \in \mathbb{Z}$ satisfying $x_1m_2m_3 + m_1x_2m_3 + m_1m_2x_3 = 1$. Then for any integers $a_1, a_2, a_3 \in \mathbb{Z}$, the general solution to the congruences $c \equiv a_i \mod m_i$ is given by

 $c \equiv a_1 x_1 m_2 m_3 + a_2 m_1 x_2 m_3 + a_3 m_1 m_2 x_3 \mod m_1 m_2 m_3.$

In our case we have $(a_1, a_2, a_3) = (3, 6, 8)$ and $(m_1, m_2, m_3) = (5, 9, 11)$. Then by inspection¹ we observe that the integers $(x_1, x_2, x_3) = (-1, 1, 1)$ satisfy the desired property:

 $x_1m_2m_3 + m_1x_2m_3 + m_1m_2x_3 = 99x_1 + 55x_2 + 45x_3 = 1.$

Therefore the general solution is

 $c \equiv a_1 x_1 m_2 m_3 + a_2 m_1 x_2 m_3 + a_3 m_1 m_2 x_3 \mod m_1 m_2 m_3$ $\equiv 3 \cdot 99 \cdot (-1) + 6 \cdot 55 \cdot 1 + 8 \cdot 45 \cdot 1 \mod 405$ $\equiv 393 \mod 405.$

Problem 2. Application of Bézout's Lemma. For any $a, b \in \mathbb{Z}$ with gcd(a, b) = 1, Bézout's Lemma tells us that ax + by = 1 for some $x, y \in \mathbb{Z}$.

(a) Prove the converse. That is, if ax + by = 1 for some $x, y \in \mathbb{Z}$, prove that gcd(a, b) = 1.

(b) Apply Bézout and part (a) to prove that

 $gcd(ab, c) = 1 \iff gcd(a, c) = 1$ and gcd(b, c) = 1.

¹It inspection didn't work we would use the matrix Euclidean algorithm.

(a): Let ax + by = 1 and $gcd(a, b) = d \ge 1$. Since d|a and d|b we have a = da' and b = db' for some $a', b' \in \mathbb{Z}$. But then we also have

$$1 = ax + by = da'x + db'y = d(a'x + b'y),$$

which since $d \ge 1$ implies that d = 1.

(b): Suppose that gcd(ab, c) = 1, so Bézout's identity implies that abx + cy = 1 for some integers $x, y \in \mathbb{Z}$. But then part (a) implies gcd(a, c) = 1 because a(bx) + c(y) = 1 and gcd(b, c) = 1 because b(ax) + c(y) = 1. Conversely, suppose that gcd(a, c) = 1 and gcd(b, c) = 1, so Bézout's identity implies that ax+cy = 1 and bx'+cy' = 1 for some integers $x, y, x', y' \in \mathbb{Z}$. But then we have

$$(ax + cy)(bx' + cy') = 1$$
$$abxx' + axcy' + cybx' + cycy' = 1$$
$$ab(xx') + c(axy' + ybx' + ycy') = 1,$$

hence from part (a) we conclude that gcd(ab, c) = 1.

Problem 3. GCD and LCM. Let $2 = p_1 < p_2 < p_3 < \cdots$ be the sequence of all primes. Then every positive integer $a \ge 2$ can be expressed in the form

$$a = p_1^{a_i} p_2^{a_2} p_3^{a_3} \cdots,$$

and is uniquely determined by the sequence of exponents a_1, a_2, a_3, \ldots

- (a) Prove that a|b if and only if $a_i \leq b_i$ for all i.
- (b) Prove that $gcd(a, b)_i = min(a_i, b_i)$ for all *i*.
- (c) Prove that $lcm(a, b)_i = max(a_i, b_i)$ for all *i*.
- (d) Combine (b) and (c) to prove that $gcd(a, b) \cdot lcm(a, b) = ab$. [Hint: $(ab)_i = a_i + b_i$.]

(a): Suppose that $a_i \leq b_i$ for all *i*, which means that $b_i = a_i + k_i$ for some non-negative integers $k_i \geq 0$. It follows that

$$b = p_1^{a_1+k_1} p^{a_2+k_2} p_3^{a_3+k_3} \dots = (p_1^{a_1} p_2^{a_2} p_3^{a_3} \dots) (p_1^{k_1} p_2^{k_2} p_3^{k_3} \dots) = a(p_1^{k_1} p_2^{k_2} p_3^{k_3} \dots),$$

and hence a|b. Conversely, suppose that a|b and consider the prime p_i . Then since $p_i^{a_i}$ divides a, it also divides b. But we know that $b = p_i^{b_i}m$ for some m satisfying $gcd(m, p_i) = 1$ and hence $gcd(m, p_i^{a_i}) = 1$. Thus we conclude from Euclid's Lemma that $p_i^{a_i}|p_i^{b_i}$, and hence $a_i \leq b_i$.

(b) and (c): For all integers $d \ge 1$ and for all primes p_i we have

$$\begin{aligned} d_i &\leqslant \gcd(a, b)_i \Leftrightarrow d| \gcd(a, b) & \text{part (a)} \\ &\Leftrightarrow d|a \text{ and } d|b \\ &\Leftrightarrow d_i \leqslant a_i \text{ and } d_i \leqslant b_i & \text{part (a)} \\ &\Leftrightarrow d_i \leqslant \min(a_i, b_i), \end{aligned}$$

which implies that $gcd(a, b)_i = min(a_i, b_i)$. Similarly, for all integers m we have

$$\begin{split} \operatorname{lcm}(a,b)_i \leqslant m_i \Leftrightarrow \operatorname{lcm}(a,b) | m & \text{part (a)} \\ \Leftrightarrow a | m \text{ and } b | m \\ \Leftrightarrow a_i \leqslant m_i \text{ and } b_i \leqslant m_i & \text{part (a)} \\ \Leftrightarrow \max(a_i,b_i) \leqslant m_i, \end{split}$$

which implies that $lcm(a, b)_i = max(a_i, b_i)$.

(d): For all integers $m, n \in \mathbb{Z}$ and for all primes p_i we note that $(mn)_i = m_i + n_i$. Furthermore, if $m_i = n_i$ for all primes p_i then we note that m = n. Thus we conclude from (b) and (c) that

$$[\gcd(a, b) \cdot \operatorname{lcm}(a, b)]_i = \gcd(a, b)_i + \operatorname{lcm}(a, b)_i$$
$$= \min(a_i, b_i) + \max(a_i, b_i)$$
$$= a_i + b_i$$
think about it
$$= (ab)_i,$$

and hence $gcd(a, b) \cdot lcm(a, b) = ab$.

Problem 4. RSA Cryptosystem. The following message has been encrypted using the RSA cryptosystem with public key (n, e) = (55, 23):

[17, 1, 33, 15, 1, 13, 20, 20, 9, 39, 26, 2, 14, 49, 13, 8, 2, 15, 1, 11]

Decrypt the message. [Hint A = 1, B = 2, C = 3, etc.]

Each message is represented by a number $0 \le m < 55$. (In this case, I only used numbers 1 through 26, corresponding to letters of the alphabet.) To encrypt the message I computed $c \equiv m^e \mod n$. To decrypt the message you should compute $m \equiv c^d \mod n$, where d is the decryption exponent.

Recall that the decryption exponent is defined by $d \equiv e^{-1} \mod (p-1)(q-1)$, where n = pq. To find d, we first factor n = 55 to obtain the primes p = 5 and q = 11. Now we need to find $d \equiv 23^{-1} \mod 40$, and we do this using the Euclidean algorithm. Each row corresponds to a true equation 23x + 40y = z:

x	y	z
0	1	40
1	0	23
-1	1	17
2	-1	6
-5	3	5
7	-4	1

We conclude that $23 \cdot 7 \equiv 40 \cdot 4 + 1 \equiv 1 \mod 40$, and hence d = 7. Finally, we raise each encrypted message c to the power of 7 mod 40. The resulting numbers are

$$[8, 1, 22, 5, 1, 7, 15, 15, 4, 19, 16, 18, 9, 14, 7, 2, 18, 5, 1, 11],$$

which translate to the following letters:

$$[h, a, v, e, a, g, o, o, d, s, p, r, i, n, g, b, r, e, a, k].$$