Problem 1. Complex Numbers as Real 2×2 **Matrices.** For any complex number $\alpha = a + bi \in \mathbb{C}$ with $a, b \in \mathbb{R}$ we define the following matrix:

$$M_{\alpha} := \begin{pmatrix} a & -b \\ b & a \end{pmatrix}$$

- (a) Check that for all $r \in \mathbb{R}$ and $\alpha \in \mathbb{C}$ we have $M_{(r\alpha)} = rM_{\alpha}$.
- (b) Check that for all $\alpha, \beta \in \mathbb{C}$ we have $M_{\alpha+\beta} = M_{\alpha} + M_{\beta}$ and $M_{\alpha\beta} = M_{\alpha}M_{\beta}$.
- (c) Check that for all $\alpha \in \mathbb{C}$ we have $\det(M_{\alpha}) = |\alpha|^2$.
- (d) Check that for all $\alpha \in \mathbb{C}$ we have $(M_{\alpha})^* = M_{(\alpha^*)}$, where the star operation denotes the transpose matrix and the complex conjugate, respectively.

(a): For all $r \in \mathbb{R}$ and $\alpha = a + bi \in \mathbb{C}$ we have

$$M_{(r\alpha)} = M_{(ra+rbi)} = \begin{pmatrix} ra & -rb\\ rb & ra \end{pmatrix} = r \begin{pmatrix} a & -b\\ b & a \end{pmatrix} = rM_{\alpha}.$$

(b): For all $\alpha = a + bi \in \mathbb{C}$ and $\beta = c + di \in \mathbb{C}$ we have

$$M_{\alpha+\beta} = M_{(a+c)+(b+d)i} = \begin{pmatrix} a+c & -(b+d) \\ b+d & a+c \end{pmatrix} = \begin{pmatrix} a & -b \\ b & a \end{pmatrix} + \begin{pmatrix} c & -d \\ d & c \end{pmatrix} = M_{\alpha} + M_{\beta}.$$

Furthermore, since $\alpha\beta = (ac - bd) + (ad + bc)i$, we have

$$M_{\alpha}M_{\beta} = \begin{pmatrix} a & -b \\ b & a \end{pmatrix} \begin{pmatrix} c & -d \\ d & c \end{pmatrix}$$
$$= \begin{pmatrix} ac - bd & -ad - bc \\ bc + ad & -bd + ac \end{pmatrix}$$
$$= \begin{pmatrix} ac - bd & -(ad + bc) \\ ad + bc & ac - bd \end{pmatrix} = M_{\alpha\beta}$$

(c): For all $\alpha = a + bi \in \mathbb{C}$ we have

$$\det(M_{\alpha}) = \det \begin{pmatrix} a & -b \\ b & a \end{pmatrix} = aa - (-b)b = a^{2} + b^{2} = |\alpha|^{2}.$$

[Remark: For all $\alpha, \beta \in \mathbb{C}$, it follows from the multiplicative property of determinants that

$$|\alpha|^2 |\beta|^2 = \det(M_\alpha) \det(M_\beta) = \det(M_\alpha M_\beta) = \det(M_{\alpha\beta}) = |\alpha\beta|^2.$$

This is another way to prove the multiplicative property of absolute value.]

[Remark: There wasn't really anything to do in this problem. I just wanted you to observe that these facts are true. In modern jargon, we say that the function $\alpha \mapsto M_{\alpha}$ is an **injective homomorphism of** \mathbb{R} -algebras.]

(d): For all $\alpha = a + bi \in \mathbb{C}$ we have

$$(M_{\alpha})^* = \begin{pmatrix} a & -b \\ b & a \end{pmatrix}^* = \begin{pmatrix} a & b \\ -b & a \end{pmatrix} = \begin{pmatrix} a & -(-b) \\ -b & a \end{pmatrix} = M_{a-bi} = M_{(\alpha^*)}.$$

Problem 2. Greatest Common Divisor. Let $a, b \in \mathbb{Z}$ with d = gcd(a, b). Since d is a common divisor of a and b we must have a = da' and b = db' for some integers $a', b' \in \mathbb{Z}$. In this case, prove that the numbers a', b' are *coprime*:

$$gcd(a',b') = 1.$$

[HInt: From Bézout's Identity we know that ax + by = d for some (non-unique) integers $x, y \in \mathbb{Z}$. Use this to show that any common divisor e|a' and e|b' must satisfy e|1.]

Proof. Let d = gcd(a, b) with a = da' and b = db' for some integers $a', b' \in \mathbb{Z}$. From Bézout's Identity there exist some $x, y \in \mathbb{Z}$ such that ax + by = d, hence we have

$$ax + by = d$$
$$da'x + db'y = d$$
$$\mathscr{A}(a'x + b'y) = \mathscr{A}$$
$$a'x + b'y = 1.$$

We will use this equation to show that gcd(a', b') = 1. To do this, let e be any common divisor of a' and b', so that a' = ea'' and b' = db'' for some integers $a'', b'' \in \mathbb{Z}$. It follows that

$$a'x + b'y = 1$$
$$ea''x + eb''y = 1$$
$$e(a''x + b''y) = 1.$$

But this implies that $e = \pm 1$, hence the greatest common divisor of a'' and b'' is 1.

Problem 3. Euclid's Lemma. For all integers $a, b, c \in \mathbb{Z}$, prove that

$$(a|bc \text{ and } gcd(a,b) = 1) \Rightarrow a|c.$$

[Hint: If gcd(a, b) = 1 then from Bézout's Identity there exist some (non-unique) integers $x, y \in \mathbb{Z}$ satisfying ax + by = 1. Multiply both sides by c to get acx + bcy = c. Now what?]

Proof. Suppose that a|bc; say ak = bc for some $k \in \mathbb{Z}$. Suppose also that gcd(a, b) = 1, hence from Bézout's Identity we have ax + by = 1 for some $x, y \in \mathbb{Z}$. Not multiply both sides by c to obtain

$$ax + by = 1$$
$$acx + bcy = c$$
$$acx + aky = c$$
$$a(cx + ky) = c.$$

We conclude that a|c, as desired.

Problem 4. Rational Root Test. Let $f(x) = c_n x^n + \cdots + c_1 x + c_0 \in \mathbb{Z}[x]$ be a polynomial of degree *n* with integer coefficients. Suppose that f(x) has a rational root $a/b \in \mathbb{Q}$ in lowest terms, i.e., with gcd(a, b) = 1. In this case, prove that we must have

$$a|c_0$$
 and $b|c_n$.

[HInt: Suppose that f(a/b) = 0. Multiply both sides by b^n and then use Euclid's Lemma.]

Proof. Let f(a/b) = 0 for some $a, b \in \mathbb{Z}$ with $b \neq 0$ and gcd(a, b) = 1. Then we have

$$f(a/b) = 0$$

$$c_n(a/b)^n + \dots + c_1(a/b) + c_0 = 0$$

$$b^n [c_n(a/b)^n + \dots + c_1(a/b) + c_0] = 0$$

$$c_n a^n + c_{n-1} a^{n-1} b + \dots + c_1 a b^{n-1} + c_0 b^n = 0.$$

By taking the term $c_0 b^n$ to one side, we have

$$c_0 b^n = -c_n a^n - c_{n-1} a^{n-1} b - \dots - c_1 a b^{n-1}$$

= $a \left[-c_n a^{n-1} - c_{n-1} a^{n-2} b - \dots - c_1 b^{n-1} \right].$

which implies that $a|c_0b^n$. Then since gcd(a,b) = 1, Euclid's Lemma implies that $a|c_0$. Similarly, by taking the term c_na^n to one side, we have

$$c_n a^n = -c_{n-1} a^{n-1} b - \dots - c_1 a b^{n-1} - c_0 b^n$$

= $b \left[-c_{n-1} a^{n-1} - \dots - c_1 a b^{n-2} - c_0 b^{n-1} \right],$

hence $b|c_n a^n$. Then since gcd(a, b) = 1, Euclid's Lemma implies that $b|c_n$.

Example. This result gives an algorithm to quickly find all of the rational roots of any polynomial with integer coefficients. For example, let $f(x) = 4x^3 - 12x^2 + 11x - 3$. If f(a/b) = 0 for some fraction $a/b \in \mathbb{Q}$ in lowest terms, then the Rational Root Test says that a|3 and b|4, which leads to a finite list of potential rational roots:

$$\frac{a}{b} \in \left\{ \pm 1, \pm 3, \pm \frac{1}{2}, \pm \frac{3}{2}, \pm \frac{1}{4}, \pm \frac{3}{4} \right\}.$$

By direct checking we find that 1, 1/2 and 3/2 are actual roots, hence

$$f(x) = 4(x-1)(x-1/2)(x-3/2).$$

This method does **not** help us to find non-rational roots.

Problem 5. The Regular 7-Gon. Let $\omega = e^{2\pi i/7}$ and $\alpha = \omega + \omega^{-1} = 2\cos(2\pi/7)$.

- (a) Combine the numbers $1, \alpha, \alpha^2, \alpha^3$ to find some polynomial $f(x) \in \mathbb{Z}[x]$ of degree 3 satisfying $f(\alpha) = 0$. [Hint: Use the fact that $\omega^3 + \omega^2 + \omega + 1 + \omega^{-1} + \omega^{-2} + \omega^{-3} = 0$.]
- (b) Use Problem 4 to show that your polynomial f(x) from part (a) has no rational roots.
- (c) Use part (b) to prove that the real number $\cos(2\pi/7) \in \mathbb{R}$ is irrational.

(a): First we compute the powers of α :

Working from outside in, we find that

$$\alpha^{3} + \alpha^{2} - 2\alpha - 1 = \omega^{3} + \omega^{2} + \omega + 1 + \omega^{-1} + \omega^{-2} + \omega^{-3} = 0.$$

Therefore we define $f(x) = x^3 + x^2 - 2x - 1 \in \mathbb{Z}[x]$ and we observe that $f(\alpha) = 0$.

(b): Suppose that f(x) has a rational root, so that f(a/b) = 0 for some integers $a, b \in \mathbb{Z}$ with gcd(a, b) = 1. Then from part (b) we must have a|1 and b|1, hence $a/b = \pm 1$. But we observe that $f(1) - 1 \neq 0$ and $f(-1) = -3 \neq 0$. Hence the polynomial f(x) has no rational roots, and it follows from part (a) that $\alpha = 2\cos(2\pi/7)$ is not rational.

(c): Assume for contradiction that $\cos(2\pi/7) = c/d$ for some integers $c, d \in \mathbb{Z}$. It follows that

$$\alpha = 2\cos\left(\frac{2\pi}{7}\right) = \frac{2c}{d} \in \mathbb{Q},$$

which contradicts part (b). Hence we conclude that $\cos(2\pi/7)$ is irrational.

[Remark: We will use this result later to prove that a regular 7-gon is not constructible with straightedge and compass.]

Problem 6. Conjugation of Complex Polynomials. For any polynomial $f(x) = \sum_{k\geq 0} a_k x^k \in \mathbb{C}[x]$ with complex coefficients, we define the *conjugate polynomial* as follows:

$$f^*(x) := \sum_{k \ge 0} a_k^* x^k.$$

(a) We can think of $\mathbb{R}[x] \subseteq \mathbb{C}[x]$ as a subring. For all $f(x) \in \mathbb{C}[x]$, prove that

$$f(x) \in \mathbb{R}[x] \quad \Leftrightarrow \quad f^*(x) = f(x).$$

- (b) For all $f(x), g(x) \in \mathbb{C}[x]$, prove $(f+g)^*(x) = f^*(x) + g^*(x)$ and $(fg)^*(x) = f^*(x)g^*(x)$.
- (c) For all $f(x) \in \mathbb{C}[x]$ use (a),(b) to prove that $f(x) + f^*(x) \in \mathbb{R}[x]$ and $f(x)f^*(x) \in \mathbb{R}[x]$.

(a): Recall that for all $a \in \mathbb{C}$ we have $a^* = a$ if and only if $a \in \mathbb{R}$. Then for all polynomials $f(x) = \sum_{k>0} a_k x^k \in \mathbb{C}[x]$ we have

$$f^*(x) = f(x) \Leftrightarrow \sum_{k \ge 0} a_k^* x^k = \sum_{k \ge 0} a_k x^k$$
$$\Leftrightarrow a_k^* = a_k \text{ for all } k \ge 0$$
$$\Leftrightarrow a_k \in \mathbb{R} \text{ for all } k \ge 0$$
$$\Leftrightarrow f(x) \in \mathbb{R}[x].$$

(b): Recall that for all $a, b \in \mathbb{C}$ we have $(a + b)^* = a^* + b^*$ and $(ab)^* = a^*b^*$. Then for all $f(x) = \sum_{k \ge 0} a_k x^k \in \mathbb{C}[x]$ and $g(x) = \sum_{k \ge 0} b_k x^k \in \mathbb{C}[x]$ we have

$$(f+g)^{*}(x) = \sum_{k \ge 0} (a_{k} + b_{k})^{*} x^{k}$$
$$= \sum_{k \ge 0} (a_{k}^{*} + b_{k}^{*}) x^{k}$$
$$= \sum_{k \ge 0} a_{k}^{*} x^{k} + \sum_{k \ge 0} b_{k}^{*} x^{k}$$
$$= f^{*}(x) + g^{*}(x)$$

and

$$(fg)^*(x) = \sum_{k \ge 0} \left(\sum_{i=1}^k a_i b_{k-i}\right)^* x^k$$
$$= \sum_{k \ge 0} \left(\sum_{i=1}^k a_i^* b_{k-i}^*\right) x^k$$
$$= \left(\sum_{k \ge 0} a_k^* x^k\right) \left(\sum_{k \ge 0} b_k^* x^k\right)$$
$$= f^*(x)g^*(x).$$

(c): For all $f(x) \in \mathbb{C}[x]$ we observe from part (b) that

$$(f+f^*)^*(x) = (f^*+f^{**})(x) = (f^*+f)(x) = (f+f^*)(x)$$

and

$$(ff^*)^*(x) = (f^*f^{**})(x) = (f^*f)(x) = (ff^*)(x).$$

Hence it follows from part (a) that $f(x) + f^*(x) \in \mathbb{R}[x]$ and $f(x)f^*(x) \in \mathbb{R}[x]$.

[Remark: We will use this last fact in our discussion of the Fundamental Theorem of Algebra. Here is a preview: Suppose that every real polynomial factors as a product of real polynomials of degrees 1 and 2. Now consider any complex polynomial $f(x) \in \mathbb{C}[x]$. Since $g(x) = f(x)f^*(x)$ has real coefficients we know that g(x) factors as a product of real polynomials of degrees 1 and 2, hence by the quadratic formula we know that g(x) splits over \mathbb{C} . Now suppose for contradiction that there exists a prime polynomial $p(x) \in \mathbb{C}[x]$ of degree ≥ 2 such that p(x)|f(x). Then we also have p(x)|g(x), which contradicts the fact that g(x) splits over \mathbb{C} . We conclude that f(x) also splits over \mathbb{C} . In summary, we have shown that the real version of the FTA implies the complex version of the FTA.]