- Base b Arithmetic. Convert the number 123456 into base b for the following values of b:
 (a) b = 2
 - (b) b = 5
 - (c) b = 16 [Use the letters A, B, C, D, E, F for 10, 11, 12, 13, 14, 15.]

I'll do them in reverse order.

(c): We divide 123456 by 16 and then repeatedly divide the quotient by 16:

$$123456 = 16 \cdot 7716 + 0$$

$$7716 = 16 \cdot 482 + 4$$

$$482 = 16 \cdot 30 + 2$$

$$30 = 16 \cdot 1 + 14$$

$$1 = 16 \cdot 0 + 1.$$

It follows that

 $123456 = 0 + 4 \cdot 16 + 2 \cdot 16^2 + 14 \cdot 16^3 + 1 \cdot 16^4.$

Since E represents 14 we express this as

 $123456 = (1E240)_{16}$.

(b): This time we divide 123456 by 5 and then divide each quotient by 5:

```
123456 = 5 \cdot 24691 + 1

24691 = 5 \cdot 4938 + 1

4938 = 5 \cdot 987 + 3

987 = 5 \cdot 179 + 2

179 = 5 \cdot 39 + 2

39 = 5 \cdot 7 + 4

7 = 5 \cdot 1 + 2

1 = 5 \cdot 0 + 1.
```

We conclude that

$$123456 = (12422311)_5$$

(a): This time I'll skip all the details:

$$123456 = (11110001001000000)_2.$$

2. Carry the One. This problem generalizes base 10 phenomena such as 2749999999 + 1 = 2750000000.

Fix a base $b \ge 2$. Then for any integers $k, r \in \mathbb{Z}$ with $k \ge 1$ prove that

$$1 + (b-1) + (b-1)b + (b-1)b^{2} + \dots + (b-1)b^{k-1} + rb^{k} = (r+1)b^{k}$$

[Hint: Use the geometric series $1 + b + \dots + b^{k-1} = (b^k - 1)/(b - 1)$.]

First we remind ourselves about the geometric series:

$$\begin{aligned} (1+b+b^2+\dots+b^{k-1})(b-1) &= (b+b^2+\dots+b^k) - (1+b+\dots+b^{k-1}) \\ &= -1+b-b+b^2-b^2+\dots+b^{k-1}-b^{k-1}+b^k \\ &= -1+0+0+\dots+0+b^k \\ &= b^k-1. \end{aligned}$$

It follows (for $b \neq 1$) that¹

$$1 + b + b^{2} + \dots + b^{k-1} = \frac{b^{k} - 1}{b - 1}$$

Now we will use this to show that

$$(\dots, r, b-1, b-1, \dots, b-1)_b + 1 = (\dots, r+1, 0, 0, \dots, 0)_b$$

(Assume that b-1 occurs k-1 times.) Indeed, the left side represents the number

$$1 + [(b-1) + (b-1)b + (b-1)b^{2} + \dots + (b-1)b^{k-1} + rb^{k} + \dots]$$

= 1 + (b - 1)(1 + b + b^{2} + \dots + b^{k-1}) + rb^{k} + \dots
= 1 + (b - 1)(b^{k} - 1)/(b - 1) + rb^{k} + \dots
= 1 + (b^{k} - 1) + rb^{k} + \dots
= b^{k} + rb^{k} + \dots
= (r + 1)b^{k} + \dots
= 0 + 0b + 0b^{2} + \dots + 0b^{k-1} + (r + 1)b^{k} + \dots

3. Lemma for the Euclidean Algorithm. Consider any positive $a, b, c, x \in \mathbb{Z}$ such that

$$a = bx + c.$$

- (a) If $d \in \mathbb{Z}$ is a common divisor of b and c, show that d also divides a.
- (b) If $d \in \mathbb{Z}$ is a common divisor of a and b, show that d also divides c.
- (c) Combine (a) and (b) to show that gcd(a, b) = gcd(b, c).

(a): Suppose that d|b and d|c, so that b = db' and c = dc' for some integers $b', c' \in \mathbb{Z}$. Since a = bx + c it follows that

$$a = bx + c$$

= $db'x + dc'$
= $d(b'x + c')$,

and hence d|a.

(b): Suppose that d|a and d|a, so that a = da' and b = db' for some integers $a', b' \in \mathbb{Z}$. Since a = bx + c it follows that

$$c = a - bx$$

= $da' - db'x$
= $d(a' - b'x)$,

and hence d|c.

¹Remark: Remind yourself what happens when |b| < 1 and k goes to infinity.

(c): We have shown that the set of common divisors of a and b is the same as the set of common divisors of b and c:

 $\{\text{common divisors of } a \text{ and } b\} = \{\text{common divisors of } b \text{ and } c\}.$

It follows that the greatest element of each set is the same, i.e., that gcd(a, b) = gcd(b, c).

4. Extended Euclidean Algorithm.

- (a) Find integers $x, y \in \mathbb{Z}$ such that 221x + 132y = 1.
- (b) Use your answer to solve the congruence $221c \equiv 7 \pmod{132}$ to find c. [Hint: From part (a) we have $221x \equiv 1 \pmod{132}$. Multiply both sides of $221c \equiv 7$ by x.]

(a): We consider the set of integer triples (x, y, r) satisfying 221x + 132y = r. Beginning with the obvious triples (1, 0, 221) and (0, 1, 132), we perform row operations until we reach a triple of the form (x, y, 1):

$$\begin{array}{c|cccc} x & y & r \\ \hline 1 & 0 & 221 \\ 0 & 1 & 132 \\ 1 & -1 & 89 \\ -1 & 2 & 43 \\ 3 & -5 & 3 \\ -43 & 72 & 1. \end{array}$$

Reminder of the method: Dividing 43 by 3 gives $43 = 14 \cdot 3 + 1$. Thus the row following (-1, 2, 43) and (3, -5, 3) is

$$(-1, 2, 43) - 14(3, -5, 3) = (-43, 72, 1).$$

We conclude that 221(-43) + 132(72) = 1. Note: This solution is **not unique**. Since 221(132k) + 132(-221k) = 0 for any k, we also have

$$221(-43+132k) + 132(72-221k) = 1 \quad \text{for any } k \in \mathbb{Z}.$$

(b): Since $132 \equiv 0 \pmod{132}$, the result from part (a) tells us that

$$1 \equiv 221(-43) + 132(72) \equiv 221(-43) + 0(72) \equiv 221(-43) \pmod{132}$$

In other words, we can kill 221 (mod 132) by multiplying by $-43 \pmod{132}$, which in standard form is 89 (mod 132). That is, we have

$$221 \cdot 89 \equiv 221 \cdot (-43) \equiv 1 \pmod{132}.$$

Thus, to solve the congruence $221c \equiv 7 \pmod{132}$ we should multiply both sides by 89:

$$221c \equiv 7$$

$$89 \cdot 221c \equiv 89 \cdot 7$$

$$1c \equiv 623$$

$$c \equiv 95 \pmod{132}.$$

This answer is unique mod 132, but it represents infinitely many integer solutions:

 $c = (any integer that is congruent to 95 \mod 132)$

= (any integer of the form 95 + 132k for some integer $k \in \mathbb{Z}$).

5. Freshman's Dream. Let $p \ge 2$ be prime.

(a) For any integer 0 < k < p, use Euclid's Lemma to prove that

$$\binom{p}{k} \equiv 0 \pmod{p}.$$

[Hint: We know that $p! = \binom{p}{k}k!(p-k)!$. Since p divides p!, Euclid's Lemma tells us that p divides $\binom{p}{k}$ or k!(p-k)! If 0 < k < p-1, show that p cannot divide k!(p-k)!.] (b) For any integers $a, b \in \mathbb{Z}$, use part (a) to prove that

$$(a+b)^p \equiv a^p + b^p \pmod{p}$$

[Hint: Use the Binomial Theorem.]

(a): Let $p \ge 2$ be prime and consider any integer 0 < k < p. The binomial coefficient $\binom{p}{k}$ satisfies the equation

$$p! = \binom{p}{k} k! (p-k)!$$
$$p(p-1)\cdots 3 \cdot 2 \cdot 1 = \binom{p}{k} k(k-1)\cdots 3 \cdot 2 \cdot 1 \cdot (p-k)(p-k-1)\cdots 3 \cdot 2 \cdot 1.$$

Since p divides the left hand side, it must also divide the right hand side:

$$p \left| \binom{p}{k} k(k-1) \cdots 3 \cdot 2 \cdot 1 \cdot (p-k)(p-k-1) \cdots 3 \cdot 2 \cdot 1 \right|$$

Since p is prime, Euclid's Lemma² tells us that p must divide one of the factors on the right hand side. However, since 0 < k < p, every factor on the right hand side is smaller than p, except for $\binom{p}{k}$. Since p cannot divide a number that is smaller than itself, we conclude that p divides $\binom{p}{k}$, which is equivalent to saying that

$$\binom{p}{k} \equiv 0 \pmod{p}.$$

(b): Let $p \ge 2$ be prime and consider any two integers $a, b \in \mathbb{Z}$. Then from part (a) and the Binomial Theorem we have

$$(a+b)^{p} \equiv a^{p} + {p \choose 1} a^{p-1}b + {p \choose 2} a^{p-2}b^{2} + \dots + {p \choose p-1} ab^{p-1} + b^{p}$$

$$\equiv a^{p} + 0a^{p-1}b + 0a^{p-2}b^{2} + \dots + 0ab^{p-1} + b^{p}$$

$$\equiv a^{p} + b^{p} \pmod{p}.$$

6. RSA Cryptosystem. You are Eve the eavesdropper. You see that Bob sent the following message to Alice using the public key (n, e) = (55, 27):

[2, 1, 33, 25, 1, 9, 4, 42, 25, 41, 1, 23, 23, 18, 17, 25, 1, 11].

Decrypt the message. [Hint: Factor n = pq as a product of primes. Then find some d such that $de \equiv 1 \pmod{(p-1)(q-1)}$; using trial and error, or using Extended Euclidean Algorithm. This is the decryption exponent. After decryption, numbers $1, \ldots, 26$ stand for letters.]

²Recall: If p is prime then Euclid's Lemma says that p|ab implies p|a or p|b.

Notice that n = 55 factors as $n = pq = 5 \cdot 11$, where p = 5 and q = 11 are prime. There, we broke the system.³ Next we need to find the decryption exponent. Recall that d satisfies

$$de + (p-1)(q-1)k = 1,$$

for some integer k whose value we don't care about. Since e = 27 and (p-1)(q-1) = 40 we want to find integers $d, k \in \mathbb{Z}$ such that

$$40k + 27d = 1$$

and this can be done with the Extended Euclidean Algorithm:

k	d	r
1	0	40
0	1	27
1	-1	13
-2	3	1.

We conclude that 27(3) + 40(-2) = 1, hence we can take d = 3 as the decryption exponent.

To encrypt a message $0 \le m < 55$, Bob computes $c = m^{27} \pmod{55}$. Then to decrypt Bob's message we compute $c^3 \pmod{55}$. The standard representative of $c^3 \pmod{55}$, i.e., the representative between 0 and 54, is guaranteed to equal m. Here is Bob's encrypted message:

[2, 1, 33, 25, 1, 9, 4, 42, 25, 41, 1, 23, 23, 18, 17, 25, 1, 11].

Raising each integer to the power of 3 and then reducing mod 55 gives

[8, 1, 22, 5, 1, 14, 9, 3, 5, 6, 1, 12, 12, 2, 18, 5, 1, 11],

which corresponds to the message

$$[h, a, v, e, a, n, i, c, e, f, a, l, l, b, r, e, a, k].$$

³If p and q were very large we would not be able to factor n = pq.