No communication devices or notes are allowed. Any student caught cheating will receive a score of zero. Each of the 4 problems is worth 6 points, for a total of 24 points.

Problem 1.

(a) Use the Extended Euclidean Algorithmn to compute the inverse of 10 mod 17.

Consider the set of triples $(k, \ell, m) \in \mathbb{Z}^3$ such that $10k + 17\ell = m$. Beginning with (0, 1, 17) and (1, 0, 10) we have the following table:

$$\begin{array}{c|ccccc} k & \ell & m \\ \hline 0 & 1 & 17 \\ 1 & 0 & 10 \\ -1 & 1 & 7 \\ 2 & -1 & 3 \\ -5 & 3 & 1 \end{array}$$

It follows that $10 \cdot (-5) = 1 \mod 17$ and hence $10^{-1} = -5 = 12 \mod 17$.

(b) Use your answer from (a) to solve the equation $10x = 8 \mod 17$.

Dividing by 10 is the same as multiplying by 12. Thus we have

$$10x = 8$$

$$x = 8 \cdot 12$$

$$x = 96$$

$$x = 11 \mod 17.$$

(c) Compute 16¹⁶ mod 17. You may assume that Fermat's Little Theorem is true.

Fermat's Little Theorem says that $a^{p-1} = 1 \mod p$ for all $a, p \in \mathbb{Z}$ with p prime and $p \nmid a$. Taking a = 16 and p = 17 gives

$$16^{16} = 1 \mod 17$$
.

Problem 2. Let \sim be an equivalence relation on a set S.

- (a) Accurately state the three axioms that \sim must satisfy.
 - (E1) $\forall x \in S, x \sim x,$
 - (E2) $\forall x, y \in S, x \sim y \Rightarrow y \sim x$,
 - (E3) $\forall x, y, z \in S, (x \sim y \land y \sim z) \Rightarrow x \sim z.$
- (b) For all $x \in S$ let $[x] := \{z \in S : x \sim z\}$. Prove that [x] = [y] implies $x \sim y$.

Assume that [x] = [y]. Then from (E1) we have $y \in [y] = [x]$, which implies that $x \sim y$.

(c) Continuing from (b), prove that $x \sim y$ implies [x] = [y].

Assume that $x \sim y$, so (E2) implies $y \sim x$. To prove that $[x] \subseteq [y]$, suppose that $z \in [x]$, which means that $x \sim z$. Then from (E3) we have $(y \sim x \land x \sim z) \Rightarrow y \sim z$ and hence $z \in [y]$. To prove that $[y] \subseteq [x]$, suppose that $z \in [y]$, which means that $y \sim z$. Then from (E3) we have $(x \sim y \land y \sim z) \Rightarrow x \sim z$ and hence $z \in [x]$.

Problem 3. Fix an integer $n \in \mathbb{Z}$ and for all integers $a, b \in \mathbb{Z}$ consider the relation

$$a \sim_n b \iff n | (a - b).$$

- (a) Prove that \sim_n is an equivalence relation on \mathbb{Z} .
 - (E1) For all $a \in \mathbb{Z}$ we have $n \cdot 0 = (a a)$, hence $n \mid (a a)$, hence $a \sim_n a$.
 - (E2) Consider $a, b \in \mathbb{Z}$ with $a \sim_n b$, so that a b = nk for some $k \in \mathbb{Z}$. Then (b-a) = n(-k) implies that n|(b-a) and hence $b \sim_n a$.
 - (E3) Consider $a, b, c \in \mathbb{Z}$ with $a \sim_n b$ and $b \sim_n c$. This means that a b = nk and $b c = n\ell$ for some $\ell \in \mathbb{Z}$. But then we have

$$a - c = (a - b) + (b - c) = nk + n\ell = n(k + \ell),$$

which implies that n|(a-c), hence $a \sim_n c$.

(b) For all $a \in \mathbb{Z}$ define the set $[a]_n = \{c \in \mathbb{Z} : a \sim_n c\}$. Prove that $[a]_n = [a']_n$ and $[b]_n = [b']_n$ imply $[a+b]_n = [a'+b']_n$.

Assume that $[a]_n = [a']_n$ and $[b]_n = [b']_n$. From Problem 2(b) and the definition of \sim_n this means that a - a' = nk and $b - b' = n\ell$ for some $k, \ell \in \mathbb{Z}$. But then we have

$$(a+b) - (a'+b') = (a-a') + (b-b') = nk + n\ell = n(k+\ell),$$

which implies that $(a+b) \sim_n (a'+b')$. Then Problem 2(c) implies $[a+b]_n = [a'+b']_n$.

(c) Continuing from (b), prove that $[a]_n = [a']_n$ and $[b]_n = [b']_n$ imply $[ab]_n = [a'b']_n$.

Assume that $[a]_n = [a']_n$ and $[b]_n = [b']_n$. From Problem 2(b) and the definition of \sim_n this means that a - a' = nk and $b - b' = n\ell$ for some $k, \ell \in \mathbb{Z}$. But then we have

$$ab - a'b' = ab - (a - nk)(b - n\ell)$$
$$= ab - (ab - an\ell - bnk + n^2k\ell)$$
$$= n(a\ell + bk - nk\ell),$$

which implies that $ab \sim_n a'b'$. Then Problem 2(c) implies $[ab]_n = [a'b']_n$.

Problem 4. Fix integers $a, b, n \in \mathbb{Z}$ and for all $k \in \mathbb{Z}$ consider the following statement:

$$P(k) = "[a^k]_n = [b^k]_n.$$

(a) Accurately state the Principle of Induction.

Let P(n) be a statement depending on an integer $n \in \mathbb{Z}$. Suppose that

 $\left\{ \begin{array}{l} \bullet \ P(b) \ \text{is true for some specific } b \in \mathbb{Z}, \ \text{and} \\ \bullet \ \text{for all } n \geq b \ \text{we have } P(n) \Rightarrow P(n+1). \end{array} \right.$

Then it follows that P(n) is true for all $n \ge b$.

(b) Assuming that P(1) is true, use induction to prove that P(k) is true for all $k \ge 1$. [Hint: You can use the result of Problem 3(c).]

Let $[a]_n = [b]_n$, i.e., P(1) is true. Now assume for induction that $[a^k]_n = [b^k]_n$, i.e., P(k) is true. Then since $[a]_n = [b]_n$ and $[a^k]_n = [b^k]_n$ it follows from Problem 3(c) that $[a \cdot a^k]_n = [b \cdot b^k]_n$ $[a^{k+1}]_n = [b^{k+1}]_n$.

In other words, P(k+1) is true. It follows by induction that P(k) is true for all $k \ge 1$.

Alternate Proof. For all $k \ge 1$ note that

$$(a^{k} - b^{k}) = (a - b)(a^{k-1} + a^{k-2}b + \dots + ab^{k-2} + b^{k-1}).$$

It follows that n|(a-b) implies $n|(a^k - b^k)$.

There are 4 problems, worth 6 points each, for a total of 24 points. This is a closed book test. Anyone caught cheating will receive a score of **zero**.

Problem 1. Hand Computations.

(a) Use Pascal's Triangle to compute the expansion of $(1+x)^5$.

Here is Pascal's Triangle:

Thus we conclude that

$$(1+x)^5 = 1 + 5x + 10x^2 + 10x^3 + 5x^4 + x^5.$$

(b) Compute the standard form of $\left[\frac{7!}{3!4!}\right]_7$.

$$\left[\frac{7!}{3!4!}\right]_7 = \left[\frac{7\cdot 6\cdot 5\cdot 4\cdot 3\cdot 2\cdot 1}{3\cdot 2\cdot 1\cdot 4\cdot 3\cdot 2\cdot 1}\right]_7 = [35]_7 = [0]_7.$$

(c) Compute the standard form of $[2^6]_7$.

$$[2^6]_7 = [2^3]_7 \cdot [2^3]_7 = [8]_7 \cdot [8]_7 = [1]_7 \cdot [1]_7 = [1]_7$$

Problem 2. Modular Arithmetic. Let $0 \neq n \in \mathbb{Z}$. Define the set $\mathbb{Z}/n := \{[a]_n : a \in \mathbb{Z}\}$ with equivalence relation $[a]_n = [b]_n \Leftrightarrow n | (a - b)$ and algebraic operations

$$[a]_n + [b]_n := [a+b]_n$$
 and $[a]_n \cdot [b]_n := [ab]_n$.

(You can assume that this is all well-defined.) Recall that \mathbb{Z}/n is a ring with additive identity element $[0]_n$ and multiplicative identity element $[1]_n$.

(a) If gcd(a, n) = 1, prove that the element $[a]_n \in \mathbb{Z}/n$ has a multiplicative inverse. [You can assume Bézout's Lemma.]

Proof. Since gcd(a, n) = 1, Bézout's Lemma says that there exist $x, y \in \mathbb{Z}$ with ax + ny = 1. Then we have

$$1 - ax = ny \implies n|(ax - 1) \implies [1]_n = [ax]_n = [a]_n \cdot [x]_n.$$

It follows that the inverse exists:

$$[a^{-1}]_n = [x]_n.$$

(b) If the element $[a]_n \in \mathbb{Z}/n$ has a multiplicative inverse, prove that there exist $x, y \in \mathbb{Z}$ with ax + ny = 1.

Proof. Suppose there exists $x \in \mathbb{Z}$ with $[a]_n \cdot [x]_n = [1]_n$. Then we have $[1]_n = [ax]_n \implies n | (1 - ax) \implies 1 - ax = ny \text{ for some } y \in \mathbb{Z}.$

(c) If there exist $x, y \in \mathbb{Z}$ with ax + ny = 1, prove that gcd(a, n) = 1.

Proof. Suppose that ax + ny = 1 for some $x, y \in \mathbb{Z}$ and let $d \in \mathbb{Z}$ be **any** common divisor of a and b, say a = dk and $b = d\ell$ for some $k, \ell \in \mathbb{Z}$. Then we have

$$1 = ax + by = (dk)x + (d\ell)y = d(kx + \ell'y),$$

which implies that $d = \pm 1$. It follows that the greatest common divisor is 1.

Problem 3. Principle of Induction.

(a) Accurately state the Principle of Induction.

Let P(n) be a statement depending on an integer $n \in \mathbb{Z}$. Suppose that

- $\left\{ \begin{array}{l} \bullet \ P(b) \ \text{is true for some specific } b \in \mathbb{Z}, \ \text{and} \\ \bullet \ \text{for all} \ n \geq b \ \text{we have} \ P(n) \Rightarrow P(n+1). \end{array} \right.$

Then it follows that P(n) is true for all $n \ge b$.

(b) For all integers $n \ge 2$ define the statement $P(n) := (1 + 2 + \dots + n) = \frac{n(n+1)}{2}$. Prove that P(2) is a true statement.

$$1+2 = \frac{2\cdot 3}{2}$$

(c) Now fix an integer $k \ge 2$ and assume for induction that P(k) is true. In this case, prove that P(k+1) is also true.

Proof. Fix an integer $k \geq 2$ and assume for induction that P(k) is true. In other words, assume that

$$1 + 2 + \dots + k = \frac{k(k+1)}{2}.$$

Then it follows that

$$1 + 2 + \dots + (k+1) = (1 + 2 + \dots + k) + (k+1)$$
$$= \frac{k(k+1)}{2} + (k+1)$$
$$= \left(\frac{k}{2} + 1\right)(k+1)$$
$$= \frac{(k+2)}{2}(k+1)$$
$$= \frac{(k+1)((k+1)+1)}{2}.$$

In other words, P(k+1) is true.

Problem 4. Binomial Theorem.

(a) Accurately state the Binomial Theorem.

For all integers $a, b, n \in \mathbb{Z}$ with $n \ge 0$ we have

$$(a+b)^n = \sum_{k=0}^n \frac{n!}{k!(n-k)!} a^k b^{n-k}.$$

(b) Let $k, p \in \mathbb{Z}$ with p prime and $1 \le k \le p-1$. In this case you can assume that p divides the integer $\frac{p!}{k!(p-k)!}$. Use this fact together with the Binomial Theorem to prove that for all $a, b \in \mathbb{Z}$ we have $[(a+b)^p]_p = [a^p + b^p]_p$.

Proof. When 1 < k < p we have assumed that

$$\left[\frac{p!}{k!(p-k)!}\right]_p = [0]_p.$$

Then using the Binomial Theorem gives

$$\begin{split} [(a+b)^{p}]_{p} &= \left[a^{p} + \sum_{k=1}^{p-1} \frac{p!}{k!(p-k)!} a^{k} b^{p-k} + b^{p}\right]_{p} \\ &= [a^{p}]_{p} + \sum_{k=1}^{p-1} \left[\frac{p!}{k!(p-k)!}\right]_{p} \cdot [a^{k} b^{p-k}]_{p} + [b^{p}]_{p} \\ &= [a^{p}]_{p} + \sum_{k=1}^{p-1} [0]_{p} \cdot [a^{k} b^{p-k}]_{p} + [b^{p}]_{p} \\ &= [a^{p}]_{p} + \sum_{k=1}^{p-1} [0]_{p} + [b^{p}]_{p} \\ &= [a^{p}]_{p} + [0]_{p} + [b^{p}]_{p} \\ &= [a^{p}]_{p} + [b^{p}]_{p}. \end{split}$$

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