

Mon Oct 7

The study of mathematics is CUMULATIVE.
Now we will see an example of this by building a major theorem from the ground up.

The theorem is the following.

Unique Factorization Theorem 2.54
(AKA. The Fundamental Theorem of Arithmetic):

Every $0 \neq n \in \mathbb{Z}$ can be expressed as

$$n = \pm p_1^{\alpha_1} p_2^{\alpha_2} \cdots p_k^{\alpha_k}$$

where p_1, p_2, \dots, p_k are primes with exponents $\alpha_1, \alpha_2, \dots, \alpha_k \in \mathbb{N}$.

And this expression is UNIQUE.

First, what is a prime?



Def: Given $p \in \mathbb{N}$ with $p > 1$,
we say that p is prime if

$$\forall a \in \mathbb{N}, a | p \Rightarrow a = 1 \text{ or } p.$$

(We say p has no nontrivial divisors)

If $n \in \mathbb{N}$ is not prime we say it is
composite. i.e.

$$\exists a \in \mathbb{Z}, a | p \text{ and } 1 < a < p.$$

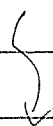
(We say n has a nontrivial divisor) //

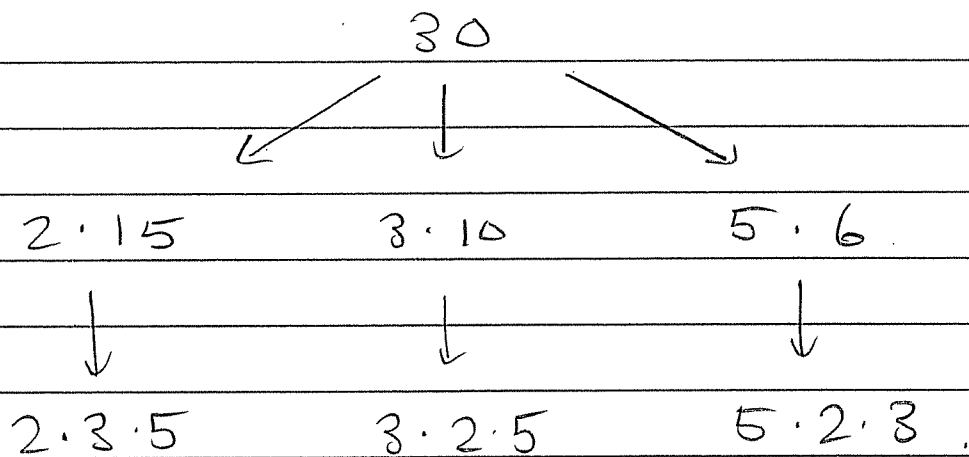
Example: The first primes are

$$2, 3, 5, 7, 11, 13, 17, 19, \dots$$

The number 30 is composite.

let's factor it!





We got the same final answer no matter how we factored 30.

That's called

UNIQUE PRIME FACTORIZATION.

$$\begin{aligned}
 30 &= 2 \cdot 3 \cdot 5 = 2 \cdot 5 \cdot 3 = 3 \cdot 2 \cdot 5 \\
 &= 3 \cdot 5 \cdot 2 = 5 \cdot 2 \cdot 3 = 5 \cdot 3 \cdot 2
 \end{aligned}$$

order doesn't matter.

Note: It didn't have to be this way.

\exists number systems without unique factorization.



Example: Consider the number system

$$\mathbb{Z}[\sqrt{-5}] := \{a + b\sqrt{-5} : a, b \in \mathbb{Z}\}.$$

Then try to factor $6 \in \mathbb{Z}[\sqrt{-5}]$:

$$\begin{array}{ccc} & 6 & \\ & \swarrow \quad \searrow & \\ 2 \cdot 3 & & (1 + \sqrt{-5})(1 - \sqrt{-5}) \\ \downarrow ? & & \downarrow ? \end{array}$$

can't be factored further!

So $6 \in \mathbb{Z}[\sqrt{-5}]$ has two different prime factorizations

Our goal is to prove that \mathbb{Z} has unique factorization

Our tool is the concept of

gcd = greatest common divisor

Def: Given integers $a, b \in \mathbb{Z}$, not both zero, consider the set of common divisors

$$\text{Div}(a, b) := \{ d \in \mathbb{Z} : d|a \text{ and } d|b \}$$

Example:

$$\text{Div}(18, 30) = \{ -6, -3, -2, -1, 1, 2, 3, 6 \}$$

We define $\text{gcd}(a, b)$ to be the highest element of $\text{Div}(a, b)$

$$\text{gcd}(a, b) := \max \text{Div}(a, b)$$

Example: $\text{gcd}(18, 30) = 6$

Q: Why does $\text{Div}(a, b)$ have a highest element?

A: Note that $\text{Div}(a, b)$ is bounded above by $\min \{ |a|, |b| \}$. Indeed, let $d \in \text{Div}(a, b)$.

}

Then $d|a \Rightarrow d \leq |d| \leq |a|$
and $d|b \Rightarrow d \leq |d| \leq |b|$

hence $d \leq \min\{|a|, |b|\}$. $///$

Since $\text{Div}(a, b)$ is bounded above, it has a highest element by Well-ordering.

This is true for subsets of \mathbb{Z} , but not subsets of \mathbb{R} .

Example: Consider the open interval

$$(0, 1) := \{x \in \mathbb{R} : 0 < x < 1\}$$

Then $(0, 1)$ is bounded above (by 1) but it has no highest element.

Proof: Assume $(0, 1)$ does have a highest element, say $x \in (0, 1)$.

But then $(1+x)/2$ is higher:

$$x < \frac{1+x}{2} < 1.$$

Contradiction. $///$

So $\sqrt{2}$ and π have no greatest common divisor.

But for all $a, b \in \mathbb{Z}$, not both zero,

$\gcd(a, b)$ exists.

Q: How to compute it?

$$\gcd(1053, 481) = ?$$

Wed Oct 9

HW 3 due Wed Oct 23

NO CLASS Fri Oct 18 (Fall Break)

NO CLASS Fri Oct 25 (I'm out of town)

Exam 1 statistics

Total	20	$A \approx 13 - 19.5$	(9)
Average	11.8	$B \approx 9.5 - 12$	(8)
Median	11	$C \approx 6 - 8.5$	(8)
St. Dev	4		

Where were we?

Given $a, b \in \mathbb{Z}$ not both 0, consider the set of common divisors

$$\text{Div}(a, b) := \{d \in \mathbb{Z} : a|d \text{ and } b|d\}$$

Then define the greatest common divisor

$$\text{gcd}(a, b) := \max \text{Div}(a, b)$$

Example:

$$\text{Div}(12, 18) = \{-6, -3, -2, -1, 1, 2, 3, \textcircled{6}\}$$

gcd

Problem: Compute $\gcd(1053, 481)$.

Two Methods.

(1) Bad Method. (Slow)

For all d from 1 to 481, test if $d|1053$ and $d|481$.
Report the largest such d .

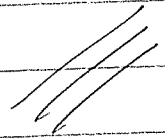
(2) Good Method (Fast).

This method depends on a trick.

★ Proposition 2.21.

Given $a, b, q, r \in \mathbb{Z}$ with $a = qb + r$
we have that

$$\gcd(a, b) = \gcd(b, r)$$



To prove this we will show an equality of sets

$$\text{Div}(a,b) = \text{Div}(b,r).$$

It will then follow that

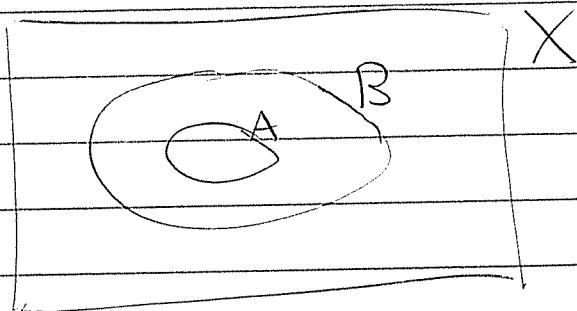
$$\begin{aligned} \text{gcd}(a,b) &= \max \text{Div}(a,b) \\ &= \max \text{Div}(b,r) = \text{gcd}(b,r). \end{aligned}$$

How can we show that two sets are equal?

Let X be a set and consider two subsets $A, B \subseteq X$.

What does it mean to say that " $A \subseteq B$ "

Picture:



Formally:

$$"A \subseteq B" \equiv " \forall x \in X, x \in A \Rightarrow x \in B "$$

To prove that " $A \subseteq B$ ":

Consider an arbitrary element $x \in A$.

Show that x is also in B .

The statement " $A = B$ " means exactly

$$"A = B" \equiv "(A \subseteq B) \text{ AND } (B \subseteq A)"$$

$$\equiv " \forall x \in X, x \in A \Leftrightarrow x \in B "$$

To prove " $A = B$ " there are two steps.

i) Prove $A \subseteq B$.

ii) Prove $B \subseteq A$.

(Just like proving \Leftrightarrow).

Proof of Prop 2.21:

Let $a = qb + r$. We will prove that $\text{Div}(a, b) = \text{Div}(b, r)$ and it follows that $\text{gcd}(a, b) = \text{gcd}(b, r)$.

First we show $\text{Div}(a, b) \subseteq \text{Div}(b, r)$.

So assume $d \in \text{Div}(a, b)$, i.e., $a = da'$ and $b = db'$ for some $a', b' \in \mathbb{Z}$.

Then we have

$$r = a - qb = da' - qdb' = d(a' - qb'),$$

hence $d|r$. It follows that $d \in \text{Div}(b, r)$.

Next we show $\text{Div}(b, r) \subseteq \text{Div}(a, b)$.

So assume $d \in \text{Div}(b, r)$, i.e., $b = db'$ and $r = dr'$ for some $b', r' \in \mathbb{Z}$.

Then we have

$$a = qb + r = qdb' + dr' = d(qb' + r'),$$

hence $d|a$. It follows that $d \in \text{Div}(a, b)$.

Let's apply the trick to compute $\gcd(1053, 481)$:

First divide 1053 by 481:

$$\begin{aligned} 1053 &= 2 \cdot 481 + 91 & \gcd(1053, 481) \\ & & = \gcd(481, 91) \end{aligned}$$

$$481 = 5 \cdot 91 + 26 \quad = \gcd(91, 26)$$

$$91 = 3 \cdot 26 + 13 \quad = \gcd(26, 13)$$

$$26 = 2 \cdot 13 + 0 \quad = \gcd(13, 0)$$

Hence

$$\gcd(1053, 481) = \gcd(13, 0) = 13$$

That was much faster than the Bad Method.

<u>Bad Method.</u>	<u>Good Method.</u>
2x481 divisions	4 divisions

This Good Method is called

"The Euclidean Algorithm"

Theorem 2.22:

To compute $\gcd(a, b)$ for $a, b \in \mathbb{Z}$ with $b \neq 0$, divide a by b . Then repeat:

$$a = q_1 b + r_1, \quad 0 \leq r_1 < |b|.$$

$$b = q_2 r_1 + r_2, \quad 0 \leq r_2 < r_1.$$

$$r_1 = q_3 r_2 + r_3, \quad 0 \leq r_3 < r_2.$$

\vdots

We obtain a sequence of remainders

$$r_0 := |b| > r_1 > r_2 > r_3 > \dots \geq 0.$$

By Well-Ordering $\exists n$ such that

$$r_n = 0 \text{ and } r_{n-1} > 0.$$

Claim: $\gcd(a, b) = r_{n-1}$

(last nonzero remainder).

Proof : By the trick (Prop 2.21) we have

$$\begin{aligned}\gcd(a, b) &= \gcd(b, r_1) \\ &= \gcd(r_1, r_2) \\ &= \gcd(r_2, r_3) \\ &\vdots\end{aligned}$$

$$= \gcd(r_{n-1}, \underbrace{r_n}_{=0}) = r_{n-1}$$



Fri Oct 11

HW 3 due Wed Oct 23

NO CLASS Fri Oct 18 (Fall Break)

NO CLASS Fri Oct 25 (I'm out of town)

Recall Prop 2.21 :

Given $a, b, q, r \in \mathbb{Z}$ with $a = qb + r$ we have

$$\gcd(a, b) = \gcd(b, r)$$

Recall The Euclidean Algorithm 2.22 :

Consider $a, b \in \mathbb{Z}$ with $b \neq 0$. To compute $\gcd(a, b)$, first divide a by b .
Then repeat.

$$a = q_1 b + r_1, \quad 0 \leq r_1 < |b|$$

$$b = q_2 r_1 + r_2, \quad 0 \leq r_2 < r_1$$

$$r_1 = q_3 r_2 + r_3, \quad 0 \leq r_3 < r_2$$

⋮

etc.

We obtain a sequence of remainders

$$|b| = r_0 > r_1 > r_2 > r_3 > \dots \geq 0.$$

By Well-Ordering, $\exists n \in \mathbb{N}$ such that

$$r_{n-1} > 0 \quad \text{and} \quad r_n = 0.$$

Claim: $\gcd(a, b) = r_{n-1}$
= last nonzero
remainder.

Proof: By Prop 2.21 we have

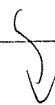
$$\begin{aligned} \gcd(a, b) &= \gcd(b, r_1) \\ &= \gcd(r_1, r_2) \\ &= \gcd(r_2, r_3) \\ &\vdots \\ &= \gcd(r_{n-1}, r_n) \\ &= \gcd(r_{n-1}, 0) = r_{n-1} \end{aligned}$$



Remark: The Euclidean Algorithm is FAST. Given $0 < a < b$, we could compute $\gcd(a, b)$ as follows:

For all d from 1 to a check if.

$$d \mid a \quad \text{and} \quad d \mid b$$



Report the largest such d . This process will require

$2 \times a$ divisions with remainder.

By contrast, computing $\gcd(a, b)$ with the Euclidean Algorithm takes

$< 2 \cdot \log_2(a)$ divisions with remainder.

(Proof omitted)

<u>Brute Force</u>	<u>Euclidean Algorithm</u>
$2a$ ☹️	$2 \cdot \log_2(a)$ 😊

Example: Compute $\gcd(21, 13)$.

It will take $< 2 \cdot \log_2(13) \approx 7.4$ steps.

$$21 = 1 \cdot 13 + 8$$

$$13 = 1 \cdot 8 + 5$$

$$8 = 1 \cdot 5 + 3$$

$$5 = 1 \cdot 3 + 2$$

$$3 = 1 \cdot 2 + \textcircled{1}$$

$$2 = 2 \cdot 1 + 0 \quad \downarrow$$

$$\gcd(21, 13) = 1$$

Definition: When $\gcd(a, b) = 1$ we say that a and b are coprime (or relatively prime).

Example: 21 & 13 are coprime.

Closely related problem:

Given $a, b, c \in \mathbb{Z}$ find all solutions $x, y \in \mathbb{Z}$ to the linear equation

$$(*) \quad \boxed{ax + by = c}$$

Note: If $d = \gcd(a, b)$, then $a = da'$ and $b = db'$ for some a', b' , hence

$$\begin{aligned} c &= ax + by \\ &= da'x + db'y \\ &= d(a'x + b'y) \end{aligned}$$

$$\implies d \mid c.$$

$\implies (*)$ has NO SOLUTION unless $\gcd(a, b) \mid c$

For simplicity, assume $c = d := \gcd(a, b)$.
Find all solutions $x, y \in \mathbb{Z}$ to
the equation

$$ax + by = d.$$

Eg. $21x + 13y = 1$.

How?

Idea: Use the Euclidean Algorithm.

Start with two easy equations

$$21 \cdot 1 + 13 \cdot 0 = 21 \quad (1)$$

$$21 \cdot 0 + 13 \cdot 1 = 13 \quad (2)$$

Combine to make more true equations

$$21 \cdot 1 + 13 \cdot (-1) = 8 \quad (3) = (1) - (2)$$

$$21(-1) + 13(2) = 5 \quad (4) = (2) - (3)$$

$$21(2) + 13(-3) = 3 \quad (5) = (3) - (4)$$

$$21(-3) + 13(5) = 2 \quad (6) = (4) - (5)$$

$$21(5) + 13(-8) = 1 \quad (7) = (5) - (6)$$

$$21(-13) + 13(21) = 0 \quad (8) = (6) - 2(7)$$

What did we find?

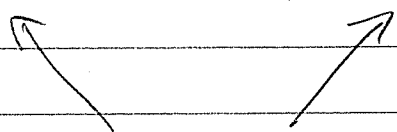
$$21(5) + 13(-8) = 1$$

is a solution ✓

Are there other solutions? Yes! Note that equation (7) + k (8) is true for all $k \in \mathbb{Z}$.

$$\begin{aligned} & [21(5) + 13(-8) = 1] \\ +k & [21(-13) + 13(21) = 0] \end{aligned}$$

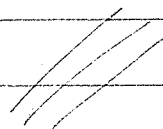
$$21(5 - 13k) + 13(-8 + 21k) = 1 + 0k$$



we get infinitely many solutions

Claim: This is the complete solution.

Proof: Postponed.



Now solve the more general problem

$$21x + 13y = c.$$

Take the solution

$$21(5 - 13k) + 13(-8 + 21k) = 1$$

and multiply both sides by c :

$$21[c(5 - 13k)] + 13[c(-8 + 21k)] = c$$

The general solution is

$$x = c(5 - 13k)$$

$$\forall k \in \mathbb{Z}$$

$$y = c(-8 + 21k)$$

★ Today's office hours 1:30 - 2:30

NO CLASS this Friday Oct 18

HW 3 due Wed Oct 23

NO CLASS next Fri Oct 25

Where were we?

Find all solutions $x, y, z \in \mathbb{Z}$ to the equation

$$3094x + 2513y = z$$

First compute $\gcd(3094, 2513)$ using the Euclidean Algorithm:

$$\begin{aligned} 3094 &= 1 \cdot 2513 + 581 & \gcd(3094, 2513) \\ 2513 &= 4 \cdot 581 + 189 & = \gcd(2513, 581) \\ 581 &= 3 \cdot 189 + 14 & = \gcd(581, 189) \\ 189 &= 13 \cdot 14 + 7 & = \gcd(189, 14) \\ 14 &= 2 \cdot 7 + 0 & = \gcd(14, 7) \\ & & = 7 \end{aligned}$$

$$\gcd(3094, 2513) = 7$$

IF $\exists x, y, z \in \mathbb{Z}$ with

$$3094x + 2513y = z.$$

Then $7 \mid \text{left side} \implies 7 \mid z$.

By contrapositive, if $7 \nmid z$ then

there is NO SOLUTION $x, y, z \in \mathbb{Z}$.

So let's try to solve

$$3094x + 2513y = 7c$$

We use the

"Extended Euclidean Algorithm" 2.25

Start with two easy equations

$$3094 \cdot 1 + 2513 \cdot 0 = 3094 \quad (1)$$

$$3094 \cdot 0 + 2513 \cdot 1 = 2513 \quad (2)$$

To save ink, just write the coefficients:

$$\begin{array}{ccc} 1 & 0 & 3094 & \textcircled{1} \\ 0 & 1 & 2513 & \textcircled{2} \end{array}$$

Now combine these to get new equations

$$\begin{array}{ccc} 1 & 0 & 3094 & \textcircled{1} \\ 0 & 1 & 2513 & \textcircled{2} \\ 1 & -1 & 581 & \textcircled{3} = \textcircled{1} - 1\textcircled{2} \\ -4 & 5 & 189 & \textcircled{4} = \textcircled{2} - 4\textcircled{3} \\ 13 & -16 & 14 & \textcircled{5} = \textcircled{3} - 3\textcircled{4} \\ -173 & 213 & 7 & \textcircled{6} = \textcircled{4} - 13\textcircled{5} \\ 359 & -442 & 0 & \textcircled{7} = \textcircled{5} - 2\textcircled{6} \end{array}$$

The last two equations are the key.

$$\begin{array}{l} 3094(-173) + 2513(213) = 7 \quad \textcircled{6} \\ 3094(359) + 2513(-442) = 0 \quad \textcircled{7} \end{array}$$

Combine them: $\textcircled{6} + k\textcircled{7}$.

$$3094(-173 + 359k) + 2513(213 - 442k) = 7 + 0 \cdot k$$

$\uparrow \qquad \qquad \qquad \uparrow$
 infinitely many solutions

Finally, multiply both sides by c :

$$3094 \underbrace{\left[c(-173 + 359k) \right]}_x + 2513 \underbrace{\left[c(213 - 442k) \right]}_y = \underbrace{7c}_z$$

This is the complete solution to

$$3094x + 2513y = z$$

$$\left. \begin{aligned} x &= c(-173 + 359k) \\ y &= c(213 - 442k) \\ z &= 7c \end{aligned} \right\} \forall k, c \in \mathbb{Z}$$


Why did we do this?

We have a theoretical reason.

Theorem (Bézout's Identity):

Let $a, b \in \mathbb{Z}$ and $d = \gcd(a, b)$. Then
 $\exists x, y \in \mathbb{Z}$ such that

$$ax + by = d.$$

Proof: We can use the Extended Euclidean Algorithm to find such $x, y \in \mathbb{Z}$.
(They are not unique.) 

Now let $a, b \in \mathbb{Z}$ (not both zero) and consider the set of \mathbb{Z} -linear combinations

$$a\mathbb{Z} + b\mathbb{Z} := \{ax + by : x, y \in \mathbb{Z}\} \subseteq \mathbb{Z}.$$

Let $d = \gcd(a, b)$ and consider the set

$$d\mathbb{Z} := \{dz : z \in \mathbb{Z}\} \subseteq \mathbb{Z}.$$

Corollary (Characterization of GCD):

$$\boxed{a\mathbb{Z} + b\mathbb{Z} = d\mathbb{Z}}$$

Proof: First show $a\mathbb{Z} + b\mathbb{Z} \subseteq d\mathbb{Z}$.

Consider any $ax + by \in a\mathbb{Z} + b\mathbb{Z}$.

Since $d|a$ and $d|b$ we have

$a = da'$ and $b = db'$ for some $a', b' \in \mathbb{Z}$.

Then

$$\begin{aligned} ax + by &= da'x + db'y \\ &= d(a'x + b'y) \in d\mathbb{Z}. \end{aligned}$$

$$\Rightarrow a\mathbb{Z} + b\mathbb{Z} \subseteq d\mathbb{Z}.$$

Next show $d\mathbb{Z} \subseteq a\mathbb{Z} + b\mathbb{Z}$.

Consider any $dz \in d\mathbb{Z}$. By Bezout's Identity $\exists x, y \in \mathbb{Z}$ such that

$$d = ax + by.$$

$$\begin{aligned} \Rightarrow dz &= (ax + by)z \\ &= a(xz) + b(yz) \in a\mathbb{Z} + b\mathbb{Z}. \end{aligned}$$

$$\Rightarrow d\mathbb{Z} \subseteq a\mathbb{Z} + b\mathbb{Z}.$$



In words:

$\gcd(a, b)$ is the smallest positive \mathbb{Z} -linear combination of a & b .

wed Oct 16

NO CLASS Friday

HW 3 due next Wed Oct 23

NO CLASS next Fri Oct 25

Recall the characterization of GCD:

Given $a, b \in \mathbb{Z}$ not both zero, consider the set of linear combinations

$$a\mathbb{Z} + b\mathbb{Z} := \{ax + by : x, y \in \mathbb{Z}\}$$

Theorem: This set has a simpler description

$$\begin{aligned} a\mathbb{Z} + b\mathbb{Z} &= \gcd(a, b)\mathbb{Z} \\ &= \{\gcd(a, b) \cdot k : k \in \mathbb{Z}\} \end{aligned}$$

Example

$$14\mathbb{Z} + 26\mathbb{Z} = 2\mathbb{Z}$$

$$= \{\dots, -6, -4, -2, 0, 2, 4, 6, \dots\}$$

Today we will consider prime factorization.

Theorem: Every $n \in \mathbb{Z}$, $n \geq 2$, has a prime factor.

Proof by Contradiction:

Suppose the statement is false and let

$$S := \{ n \geq 2 : n \text{ has no prime factor} \}$$

By assumption $S \neq \emptyset$ hence by Well-ordering \exists smallest element, say $m \in S$.

[Remark: m is a "minimal criminal"]

Since $m \in S$, m is not prime (since otherwise m is a prime factor of itself).

By definition this means $\exists a, b \in \mathbb{Z}$ such that

$$m = ab \quad \text{and} \quad 1 < a, b < m.$$

Since $a < m \implies a \notin S$, we know that a does have a prime factor,

say $a = pk$ where p is prime. But then

$$m = ab = pkb = p(kb)$$

$\Rightarrow m$ has a prime factor, contradicting the fact that $m \in S$

Theorem: Every $n \in \mathbb{N}$, $n \geq 2$, can be written as a product of primes

Proof by contradiction:

Assume the statement is false. Then by Well-Ordering there exists a smallest $m \geq 2$ with no prime factorization.


This m is not prime, but it has a prime factor by the previous Theorem, say $m = pk$ where $1 < k < m$.

Since m was minimal we know that k has a prime factorization, say

$$k = p_1 p_2 \cdots p_r$$

But then m has a prime factorization

$$m = p_1 p_2 \cdots p_r.$$

Contradiction. 

Corollary: Every $n \in \mathbb{Z}$, $n \neq 0$, can be written as a product of primes (times ± 1).

Example:

- 1 is not prime, but it is 1 times a product of primes (i.e. no primes)
- if $n < 0$ then $-n > 0$ is 1 times a product of primes.

$$n = (-1) \cdot \text{product of primes}$$

- $$\begin{aligned} 364 &= 2 \cdot 2 \cdot 7 \cdot 13 \\ &= 2 \cdot (-2) \cdot (-7) \cdot 13 \\ &= (-2) \cdot (-2) \cdot (-7) \cdot (-13) \end{aligned}$$

etc.

$$\begin{aligned} \circ -8 &= (-1) \cdot 2 \cdot 2 \cdot 2 \\ &= (-2)(-2)(-2) \\ &= 2(-2)2 \\ &\text{etc.} \end{aligned}$$

Don't worry about ± 1 !

We know prime factorizations EXIST.
Are they UNIQUE?

We need one final tool:

Theorem 2.53 ("Euclid's Lemma"):

Let p be prime. Then $\forall a, b \in \mathbb{Z}$,

$$p \mid ab \implies p \mid a \text{ OR } p \mid b$$

Proof: We will prove the equivalent statement:

$$p \mid ab \text{ AND } p \nmid a \implies p \mid b$$

So assume $p \mid ab$, say $ab = pk$,
and assume that $p \nmid a$. Then we
have $\gcd(a, p) = 1$ (Why?)

By Bézout's Identity $\exists x, y \in \mathbb{Z}$ such
that $1 = ax + py$. Multiplying
both sides by b gives

$$b = abx + pby.$$

$$= pkx + pby$$

$$= p(kx + by).$$

Hence $p \mid b$, as desired



Multiplying both sides of $1 = ax + py$ by
 b was a "trick" (an "uphill step" in
the proof), and it worked!

Mon Oct 21

HW 3 due wed
No CLASS Friday

Recap:

• $\gcd(a, b) := \max \text{Div}(a, b)$, where

$$\text{Div}(a, b) := \{ d \in \mathbb{Z} : d|a \text{ AND } d|b \}$$

• If $a = qb + r$ with $b \neq 0$ then

$$\text{Div}(a, b) = \text{Div}(b, r)$$

$$\implies \gcd(a, b) = \gcd(b, r)$$

• This + Division Algorithm. = The Euclidean Algorithm. It allows us to compute $\gcd(a, b)$ quickly.

• Moreover, the Extended Euclidean Algorithm allows us to solve the equation

$$ax + by = \gcd(a, b)$$

for $x, y \in \mathbb{Z}$.

- Then we can characterize the GCD in a nice way:

$$a\mathbb{Z} + b\mathbb{Z} = \gcd(a,b)\mathbb{Z}$$

- We say $p \in \mathbb{Z}$ is prime if $\forall a, b \in \mathbb{Z}$

$$p = ab \implies a = \pm 1 \text{ OR } b = \pm 1$$

- We say $n \in \mathbb{Z}$ is composite if it is not prime, i.e. $\exists a, b \in \mathbb{Z}$ with

$$a, b \neq \pm 1 \text{ AND } n = ab.$$

[We say these a, b are proper factors of n (or nontrivial factors)]

- We used well-ordering to show that $\forall n \in \mathbb{Z}, n \neq 0$, we have an expression

$$n = \pm p_1^{e_1} p_2^{e_2} \cdots p_k^{e_k}$$

where p_1, p_2, \dots, p_k are prime.

Q: Is this expression unique?

A: How could it not be?

Here is the key point:

★ Theorem 2.53 ("Euclid's Lemma")
(Book VII Prop 30 in Euclid) :

Let $p \in \mathbb{Z}$ be prime. Then $\forall a, b \in \mathbb{Z}$
we have

$$\boxed{p \mid ab \implies p \mid a \text{ OR } p \mid b.}$$

[Note: This fails for non-prime p .

Eg. $p=4$. We have

$$4 \mid 12 = 2 \cdot 6 \text{ BUT } 4 \nmid 2 \text{ AND } 4 \nmid 6.]$$

We will prove the equivalent statement

$$p \mid ab \text{ AND } p \nmid a \implies p \mid b.$$

[Recall:

$$"P \implies (Q \text{ OR } R)" \equiv "(P \text{ AND NOT } Q) \implies R"]$$

Proof: Let $p \in \mathbb{Z}$ be prime. Suppose that $p \mid ab$, say $ab = pk$, and suppose that $p \nmid a$.

Let $d = \gcd(a, p)$. Since $d \mid p$ we have $d = 1$ or $d = p$. But since $d \mid a$ and $p \nmid a$ we conclude $d = 1$.

By Bézout's identity $\exists x, y \in \mathbb{Z}$ such that

$$1 = ax + py.$$

Now multiply both sides by b to get

$$\begin{aligned} b &= b(ax + py) \\ &= bax + bpy \\ &= abx + pby \\ &= pkx + pby \\ &= p(kx + by). \end{aligned}$$

Hence $p \mid b$, as desired.



More Generally: Let p be prime.

If $p \mid a_1 a_2 \cdots a_k$ then $p \mid a_i$ for some i .

[How would you prove this?] _o

Now we have everything we need to prove unique prime factorization.

★ Unique Factorization Theorem 2.54
("Fundamental Theorem of Arithmetic").

Every integer $n \neq 0$ has a unique expression as a product of primes (apart from reordering and negative signs).

Proof Idea: Suppose n can be expressed

$$n = p_1 p_2 \cdots p_k = q_1 q_2 \cdots q_l,$$

where $p_1, \dots, p_k, q_1, \dots, q_l$ are prime.

Since $p_1 \mid n$ we have

$$p_1 \mid q_1 q_2 \cdots q_l$$

Since p_i is prime, Euclid's Lemma says that we have $p_i \mid q_i$ for some i .
Since p_i and q_i are both prime, this implies $p_i = \pm q_i$

Cancel the common factor from both sides to get

$$p_2 p_3 \cdots p_k = \pm (q_1 \cdots q_{i-1} q_{i+1} \cdots q_l)$$

Now repeat until you're done...

Is that a proof? NOT REALLY

Formal Proof:

Suppose \exists integer $\neq 0$ with two different prime factorizations. By Well-Ordering, let m be the smallest such integer and write

$$m = p_1 p_2 \cdots p_k = q_1 q_2 \cdots q_l \quad (*)$$

Different in some nontrivial way.

Since $p_1 \mid m$ we have $p_1 \mid q_1 q_2 \cdots q_l$ and
Euclid's Lemma $\Rightarrow p_1 \mid q_i$ for some i .
Since p_1, q_i are prime this implies
that $p_1 = \pm q_i$

Cancel p_1 from $(*)$ to get

$$m' = p_2 p_3 \cdots p_k = \pm q_1 \cdots q_{i-1} q_{i+1} \cdots q_l$$

still different

with $|m'| < |m|$. But this contradicts
the fact that m was minimal



[Remarks:

- We skipped some details, didn't we?
- Did we skip anything important?
- Would adding details make the proof better, or worse?



Wed Oct 23

HW 3 is due NOW.

NO CLASS Friday

Let's look at HW 3 Problem 3:

What is a fraction?

Given $a, b \in \mathbb{Z}$ with $b \neq 0$ we define an abstract symbol:

$$\text{" } \frac{a}{b} \text{"}$$

We declare rules for "multiplying" and "adding" abstract symbols:

$$\text{" } \frac{a}{b} \text{"} \cdot \text{" } \frac{c}{d} \text{"} := \text{" } \frac{ac}{bd} \text{"}$$

$$\text{" } \frac{a}{b} \text{"} + \text{" } \frac{c}{d} \text{"} := \text{" } \frac{ad+bc}{bd} \text{"}$$

We declare when two abstract symbols are "equal":

$$\frac{a}{b} = \frac{c}{d} \iff ad = bc.$$

Definition: The set

$$\mathbb{Q} := \left\{ \frac{a}{b} : a, b \in \mathbb{Z}, b \neq 0 \right\}$$

is called the system of rational numbers.

We can think of \mathbb{Z} as a subset of \mathbb{Q} by identifying

$$n \in \mathbb{Z} \iff \frac{n}{1} \in \mathbb{Q}$$

The benefit of \mathbb{Q} is that we can divide by nonzero elements.

If $x = \frac{a}{b} \neq 0$, then $a \neq 0$, hence the symbol $\frac{b}{a}$ exists and we have

$$\frac{a}{b} \cdot \frac{b}{a} = 1.$$

Every nonzero rational number has a multiplicative inverse.

Note that fractions do not have a unique representation:

$$-\frac{3}{4} = \frac{3}{-4} = \frac{6}{-8} = -\frac{6}{8} = \text{etc.}\dots$$

Q: Does each $x \in \mathbb{Q}$ have a best representation?

We will say that $x = a/b$ is in lowest terms if

- $b > 0$
- $\gcd(a, b) = 1$

Theorem: Every $x \in \mathbb{Q}$ can be represented uniquely in lowest terms.

[Remark: We have already used this when we proved that $\sqrt{2} \notin \mathbb{Q}$.
But we never proved it.]

Proof of Existence (HW 3.3):

Let $x = a/b \in \mathbb{Q}$. We can assume that $b > 0$, otherwise we just write $a/b = (-a)/(-b)$.

[Note: $\frac{a}{b} = \frac{-a}{-b}$ because $a(-b) = (-a)b$.]

Now suppose that $d = \gcd(a, b)$, with $a = da'$ and $b = db'$. Then we have

$$x = \frac{a}{b} = \frac{da'}{db'} = \frac{a'}{b'}$$

because $da'b' = db'a'$. Since $b > 0$ and $d > 0$ we have $b' > 0$.

We claim that $\gcd(a', b') = 1$.

Indeed, by Bézout's Identity

$\exists x, y \in \mathbb{Z}$ such that

$$d = ax + by.$$

$$d = da'x + db'y$$

$$d = d(a'x + b'y)$$

$$1 = a'x + b'y.$$

Now any common divisor of a' and b' must divide 1. [Suppose $a' = ka''$ and $b' = kb''$. Then

$$\begin{aligned} 1 &= a'x + b'y \\ &= ka''x + kb''y \\ &= k(a''x + b''y) \end{aligned} \quad]$$

Hence $\gcd(a', b') = 1$ and $x = a'/b'$ is in lowest terms.

Proof of Uniqueness:

Suppose that $x = \frac{a_1}{b_1} = \frac{a_2}{b_2}$ with

- $b_1 > 0, b_2 > 0$
- $\gcd(a_1, b_1) = \gcd(a_2, b_2) = 1$.

We want to show that $a_1 = a_2$ and $b_1 = b_2$.

Since $a_1/b_1 = a_2/b_2$ we have

$$a_1 b_2 = a_2 b_1.$$

We claim that $b_1 \mid b_2$.

Indeed, since $\gcd(a_2, b_2) = 1 \exists x, y \in \mathbb{Z}$
such that

$$1 = a_2 x + b_2 y$$

Multiply both sides by b_1 to get

$$\begin{aligned} b_1 &= a_2 b_1 x + b_2 b_1 y \\ &= a_1 b_2 x + b_2 b_1 y \\ &= b_2 (a_1 x + b_1 y). \end{aligned}$$

Hence $b_1 \mid b_2$. And similarly $b_2 \mid b_1$.

So we have $b_1 = s b_2$ and $b_2 = t b_1$, for
some $s, t \in \mathbb{Z}$. Hence

$$b_1 = s b_2 = s t b_1$$

$$\Rightarrow b_1 - s t b_1 = 0$$

$$\Rightarrow (1 - s t) b_1 = 0$$

Since $b_1 \neq 0$ we get $1 - s t = 0$,
or $s t = 1$. We conclude that
 $s, t = \pm 1$, hence $b_1 = \pm b_2$



Since $b_1 > 0$ and $b_2 > 0$ by assumption,
we get $b_1 = b_2$, and

$$a_1 \cancel{b_1} = a_2 b_2 = a_2 \cancel{b_1} \implies a_1 = a_2$$

