Math 230 D	Fall 2012
Homework 6	Drew Armstrong

Problem 1. For each integer $n \ge 0$, let P(n) be the statement: "any set of size n has 2^n subsets." Use induction to prove that P(n) is true for all $n \ge 0$. [Hint: Let A be an arbitrary set of size n and let $x \in A$ be some fixed element. Then every subset of A either contains x or does not. How many subsets are there of each type? [Hint: By induction, there are 2^{n-1} subsets of A that do **not** contain x, since these are just the subsets of $A \setminus \{x\}$. Show that there are also 2^{n-1} subsets that **do** contain x.]]

Proof. We wish to show that P(n) = T for all $n \ge 0$. First note that P(0) = T because there is only one set of size 0 — the empty set \emptyset — and it has exactly $2^0 = 1$ subset — itself. Now fix an arbitrary $k \ge 0$ and (OPEN MENTAL PARENTHESIS. suppose that P(k) = T. In this case we wish to show that P(k + 1) = T. So let A be an arbitrary set of size k + 1 and fix some element $x \in A$. Each subset of A either contains x or does not. The subsets of Athat do **not** contain x are precisely the subsets of $A \setminus \{x\}$, and since P(k) = T we know there are 2^k of these. On the other hand, for each subset of A that does not contain x there is a unique subset that does — namely, we just add x. Hence there are **also** 2^k of these. We conclude that A has $2^k + 2^k = 2 \cdot 2^k = 2^{k+1}$ subsets, hence P(k + 1) = T. CLOSE MENTAL PARENTHESIS.) We have shown that P(0) = T and $P(k) \Rightarrow P(k + 1)$ for all $k \ge 0$. By induction we conclude that P(n) = T for all $n \ge 0$.

Problem 2.

- (a) Let $a, b, c \in \mathbb{Z}$ with gcd(a, b) = 1. If a|c and b|c, prove that ab|c. [Hint: Use Bézout to write ax + by = 1 and multiply both sides by c.]
- (b) In class we proved *Fermat's little Theorem*, which says that if $p \in \mathbb{Z}$ is prime and gcd(a, p) = 1 (i.e. if p doesn't divide a), then we have $a^{p-1} = 1 \mod p$. To apply this to cryptography we need a slightly more general result:

Given integers $a, p, q \in \mathbb{Z}$ with p and q prime and with gcd(a, pq) = 1 (i.e. with $p \not\mid a$ and $q \not\mid a$), we have $a^{(p-1)(q-1)} = 1 \mod pq$.

Prove this result. [Hint: You may assume Fermat's little Theorem. First prove that q divides $a^{(p-1)(q-1)} - 1$. The same argument works for p. Then use part (a).]

Proof. To prove (a) suppose that a|c and b|c (say c = ak and $c = b\ell$) with a, b coprime. By Bézout, there exist integers $x, y \in \mathbb{Z}$ such that ax + by = 1. Multiply both sides by c to get

$$ax + by = 1$$
$$(ax + by)c = c$$
$$axc + byc = c$$
$$axb\ell + byak = c$$
$$ab(x\ell + yk) = c,$$

hence ab|c. To prove (b) consider $a, p, q \in \mathbb{Z}$ with p and q prime and with gcd(a, pq) = 1(i.e. with $p \not| a$ and $q \not| a$). We wish to show that pq divides $a^{(p-1)(q-1)} - 1$. To see this, first note that q does **not** divide a^{p-1} since if it did then q would also divide a (by the Extended Euclid's Lemma HW5.2), contradiction. Hence by Fermat's little Theorem we conclude that q divides $(a^{p-1})^{q-1} - 1 = a^{(p-1)(q-1)} - 1$. Similarly we see that p divides $a^{(p-1)(q-1)} - 1$. Then since p and q are coprime (indeed, they are both prime), part (a) implies that pq divides $a^{(p-1)(q-1)} - 1$, as desired. [I agree, it doesn't seem that this should be the foundation of modern cryptography, but it is.]

Problem 3. Use the Binomial Theorem to prove the following:

- (a) $\binom{n}{0} + \binom{n}{1} + \binom{n}{2} + \dots + \binom{n}{n} = 2^n$ for all $n \ge 1$.
- (b) $\binom{n}{0} \binom{n}{1} + \binom{n}{2} \dots + (-1)^n \binom{n}{n} = 0$ for all $n \ge 1$.
- (c) $0\binom{n}{0} + 1\binom{n}{1} + 2\binom{n}{2} + \dots + n\binom{n}{n} = n2^{n-1}$ for all $n \ge 1$.

[Hint: The proofs are one-liners. What is the derivative $\frac{d}{dx}$ of $(1+x)^n$?]

Proof. Recall the Binomial Theorem:

(1)
$$(1+x)^n = \binom{n}{0} + \binom{n}{1}x + \binom{n}{2}x^2 + \dots + \binom{n}{n}x^n.$$

Since this is an equation of *polynomials*, it remains true if we substitute any value for x. Putting x = 1 in equation (1) yields

$$2^{n} = \binom{n}{0} + \binom{n}{1} + \binom{n}{2} + \dots + \binom{n}{n},$$

and putting x = -1 in equation (1) yields

$$0 = 0^{n} = {\binom{n}{0}} - {\binom{n}{1}} + {\binom{n}{2}} - \dots + (-1)^{n} {\binom{n}{n}}.$$

We may also differentiate equation (1) by x to get another equation of polynomials:

(2)
$$n(1+x)^{n-1} = 0\binom{n}{0} + 1\binom{n}{1} + 2\binom{n}{2}x + \dots + n\binom{n}{n}x^{n-1}$$

Then putting x = 1 in equation (2) yields

$$n2^{n-1} = 0\binom{n}{0} + 1\binom{n}{1} + 2\binom{n}{2} + \dots + n\binom{n}{n}.$$

Problem 4. Note that we can write

$$\binom{n}{k} = \frac{n!}{k!(n-k)!} = \frac{(n)_k}{k!}$$

where $(n)_k := n(n-1)\cdots(n-(k-1))$. Why would we do this? Because the expression $(z)_k$ makes sense for any positive integer k and any complex number $z \in \mathbb{C}$. Thus we can define $\binom{z}{k} := (z)_k/k!$ for any $k \in \mathbb{N}$ and $z \in \mathbb{C}$. Prove that for all $n, k \in \mathbb{N}$ we have

$$\binom{-n}{k} = (-1)^k \binom{n+k-1}{k}.$$

Proof. Let $n, k \in \mathbb{N}$ be positive integers. Then we have

$$\binom{-n}{k} = \frac{(-n)_k}{k!} = \frac{(-n)(-n-1)(-n-2)\cdots(-n-(k-1))}{k!}$$

$$= \frac{(-1)^k (n)(n+1)(n+2)\cdots(n+(k-1))}{k!}$$

$$= (-1)^k \frac{(n+k-1)(n+k-2)\cdots(n+2)(n+1)(n)}{k!}$$

$$= (-1)^k \frac{(n+k-1)_k}{k!} = (-1)^k \binom{n+k-1}{k}.$$

[One could alternatively show that $\binom{-n}{k}$ satisfies the correct recurrence and initial conditions.]

Problem 5. Let $x, z \in \mathbb{C}$ be complex numbers with |x| < 1. Newton's Binomial Theorem says that

$$(1+x)^{z} = 1 + {\binom{z}{1}}x + {\binom{z}{2}}x^{2} + {\binom{z}{3}}x^{3} + \cdots$$

where the right hand side is a convergent infinite series. Use this to obtain an infinite series expansion of $(1 + x)^{-2}$ when |x| < 1. [Hint: Apply Problem 4.]

By Problem 4, we know that

$$\binom{-2}{k} = (-1)^k \binom{2+k-1}{k} = (-1)^k \binom{k+1}{k} = (-1)^k (k+1)$$

for all $k \in \mathbb{N}$. Then for any $x \in \mathbb{C}$ with |x| < 1, Newton tells us that

$$\frac{1}{(1+x)^2} = 1 - 2x + 3x^2 - 4x^3 + \dots = \sum_{k \ge 0} (-1)^k (k+1)x^k,$$

where the infinite series on the right is convergent. We could alternatively get this by differentiating the well-known geometric series

$$\frac{1}{1+x} = 1 - x + x^2 - x^3 + x^4 - \cdots$$

What do you get if you integrate the geometric series? Answer:

$$\log(1+x) = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \cdots$$

The geometric series is useful.